PRELIMINARY

National Semiconductor

COP888CG Single-Chip microCMOS Microcontroller

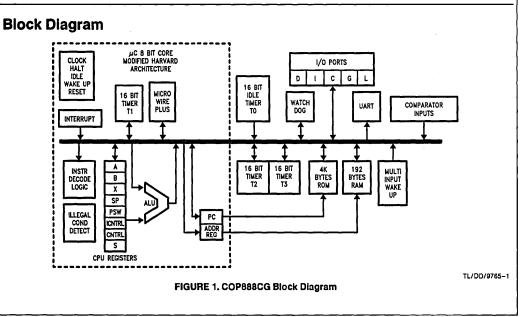
General Description

The COP888 family of microcontrollers uses an 8-bit single chip core architecture fabricated with National Semiconductor's M²CMOS™ process technology. The COP888CG is a member of this expandable 8-bit core processor family of microcontrollers. (Continued)

Features

- Low cost 8-bit microcontroller
- Fully static CMOS, with low current drain
- Two power saving modes: HALT and IDLE
- 1 µs instruction cycle time
- 4096 bytes on-board ROM
- 192 bytes on-board RAM
- Single supply operation: 2.5V-6V
- Full duplex UART
- Two analog comparators
- MICROWIRE/PLUS™ serial I/O
- WatchDog and Clock Monitor logic
- Idle Timer
- Multi-Input Wakeup (MIWU) with optional interrupts (8)
- Fourteen multi-source vectored interrupts servicing
 External Interrupt
 - Idle Timer T0
 - Three Timers (Each with 2 Interrupts)
 - MICROWIRE/PLUS
 - Multi-Input Wake Up
 - Software Trap
 - UART (2)
 - Default VIS

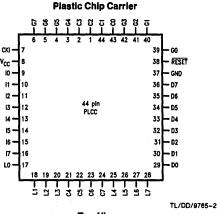
- Three 16-bit timers, each with two 16-bit registers supporting:
 - Processor Independent PWM mode
 - External Event counter mode
 - Input Capture mode
- 8-bit Stack Pointer SP (stack in RAM)
- Two 8-bit Register Indirect Data Memory Pointers (B and X)
- Versatile instruction set
- True bit manipulation
- Memory mapped I/O
- BCD arithmetic instructions
- Package: 44 PCC or 40 N or 28 N or 28 PCC — 44 PCC with 39 I/O pins
 - 40 N with 35 I/O pins
 - 28 PCC or 28 N, each with 23 I/O pins
- Software selectable I/O options
 TRI-STATE® Output
 - TRI-STATE® Outpu
 - Push-Pull Output - Weak Pull Up Input
 - High Impedance Input
- Schmitt trigger inputs on ports G and L
- Temperature ranges: -40°C to +85°C,
 - -55°C to +125°C
- ROMless mode for accurate emulation and external program memory capability
- Single chip COP888CGP piggy back emulation device
- Real time emulation and full program debug offered by National's Development Systems



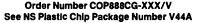
General Description (Continued)

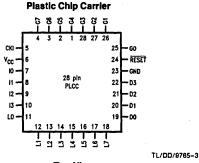
It is a fully static part, fabricated using double-metal silicon gate microCMOS technology. Features include an 8-bit memory mapped architecture, MICROWIRE/PLUS serial I/O, three 16-bit timer/counters supporting three modes (Processor Independent PWM generation, External Event counter, and Input Capture mode capabilities), full duplex UART, two comparators, and two power savings modes (HALT and IDLE), both with a multi-sourced wakeup/interrupt capability. This multi-sourced interrupt capability may

Connection Diagrams



Top View



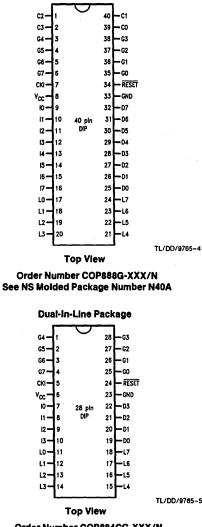


Top View

Order Number COP884CG-XXX/V See NS Plastic Chip Package Number V28A

also be used independent of the HALT or IDLE modes. Each I/O pin has software selectable configurations. The COP888CG operates over a voltage range of 2.5V to 6V. High throughput is achieved with an efficient, regular instruction set operating at a maximum of 1 μ s per instruction rate. The COP888CG may be operated in the ROMIess mode to provide for accurate emulation and for applications requiring external program memory.





Order Number COP884CG-XXX/N See NS Molded Package Number N28A

FIGURE 2a. COP888CG Connection Diagrams

Connection Diagrams (Continued)

		888CG Pinouts for				
Port	Туре	Alt. Fun	Alt. Fun	28-Pin Pack.	40-Pin Pack.	44-Pin Pack.
LO	1/0	MIWU		11	17	17
L1	1/0	MIWU	СКХ	12	18	18
L2	1/0	MIWU	TDX	13	19	19
L3	1/0	MIWU	RDX	14	20	20
L4	1/0	MIWU	T2A	15	21	25
L5	1/0	MIWU	T2B	16	22	26
L6	1/0	MIWU	T3A	17	23	27
L7	1/0	MIWU	T3B	18	24	28
G0	I/O	INT		25	35	39
G1	WDOUT			26	36	40
G2	1/0	T1B		27	37	41
G3	1/0	T1A		28	38	42
G4	1/0	SO		1	3	3
G5	1/0	SK		2	4	4
G6	1	SI		3	5	5
G7	I/CKO	HALT Restart		4	6	6
D0	0	ROM DATA*		19	25	29
D1	0	PCL*		20	26	30
D2	0	EMUL*		21	27	31
D3	0	PCU*		22	28	32
10	1			7	9	9
11	I	COMP1IN-		8	10	10
12	i	COMP1IN+		9	11	11
13	1	COMP1OUT		10	12	12
14	I.	COMP2IN-			13	13
15	I	COMP2IN+			14	14
16	I	COMP2OUT			15	15
17	1				16	16
D4	0	S CLOCK*	0		29	33
D5	0	HALTSEL*			30	34
D6	0	LOAD*			31	35
D7	0	D DATA•			32	36
C0	I/O				39	43
C1	I/O				40	44
C2	ı/o				1	1
Сз	I/O				2	2
C4	I/O					21
C5	I/O		1			22
C6	I/O					23
C7	1/0					24
Vcc				6	8	8
GND				23	33	37
СКІ				5	7	7
RESET				24	34	38

COP888CG Pinouts for 28-, 40- and 44-Pin Packages

• = Only in the ROMless Mode

Absolute Maximum Ratings

If Military/Aerospace specified devices are required, please contact the National Semiconductor Sales Office/Distributors for availability and specifications.

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Supply Voltage (V _{CC})	7V			
Voltage at Any Pin	-0.3V to V _{CC} + 0.3V			
ESD Susceptibility (Note 4)	2000V			
Total Current into V _{CC} Pin (Source)	100 mA			

Total Current out of GND Pin (Sink)

Storage Temperature Range

110 mA -65°C to +140°C

Note: Absolute maximum ratings indicate limits beyond which damage to the device may occur. DC and AC electrical specifications are not ensured when operating the device at absolute maximum ratings.

DC Electrical Characteristics $-40^{\circ}C \le T_A \le +85^{\circ}C$ unless otherwise specified

Parameter	Conditions	Min	Тур	Max	Units
Operating Voltage		2.5		6	v
Power Supply Ripple (Note 1)	Peak-to-Peak			0.1 V _{CC}	v
Supply Current (Note 2) CKI = 10 MHz CKI = 4 MHz	$V_{CC} = 6V, t_c = 1 \ \mu s$ $V_{CC} = 2.5V, t_c = 2.5 \ \mu s$			15 2	mA mA
HALT Current (Note 3)	$V_{CC} = 6V, CKI = 0 MHz$		<1		μA
IDLE Current CKI = 10 MHz CKI = 4 MHz	$V_{CC} = 6V, t_c = 1 \ \mu s$ $V_{CC} = 2.5V, t_c = 2.5 \ \mu s$			5 0.6	mA mA
Input Levels RESET Logic High Logic Low CKI (External and Crystal Osc. Modes)		0.8 V _{CC}		0.2 V _{CC}	v v
Logic High Logic Low All Other Inputs		0.7 V _{CC}		0.2 V _{CC}	v v
Logic High Logic Low		0.7 V _{CC}		0.2 V _{CC}	v v
Hi-Z Input Leakage	$V_{CC} = 6V$	-2		+2	μΑ
Input Pullup Current	$V_{CC} = 6V$	40		250	μΑ
G and L Port Input Hysteresis			0.05 V _{CC}		v
Output Current Levels D Outputs Source	V _{CC} = 4V, V _{OH} = 3.3V	0.4			mA
Sink	$V_{CC} = 4V, V_{OH} = 3.5V$ $V_{CC} = 2.5V, V_{OH} = 1.8V$ $V_{CC} = 4V, V_{OL} = 1V$ $V_{CC} = 2.5V, V_{OL} = 0.4V$	0.4 0.2 10 2.0			mA mA mA
All Others					
Source (Weak Pull-Up Mode)	$V_{CC} = 4V, V_{OH} = 2.7V$ $V_{CC} = 2.5V, V_{OH} = 1.8V$	10 2.5		100 33	μΑ μΑ
Source (Push-Pull Mode)	$V_{CC} = 4V, V_{OH} = 3.3V$ $V_{CC} = 2.5V, V_{OH} = 1.8V$	0.4 0.2			mA mA
Sink (Push-Pull Mode)	$V_{CC} = 4V, V_{OL} = 0.4V$ $V_{CC} = 2.5V, V_{OL} = 0.4V$	1.6			mA mA
TRI-STATE Leakage		-2		+2	μΑ

Note 1: Rate of voltage change must be less then 0.5 V/ms.

Note 2: Supply current is measured after running 2000 cycles with a square wave CKI input, CKO open, inputs at rails and outputs open.

Note 3: The HALT mode will stop CKI from oscillating in the RC and the Crystal configurations. Test conditions: All Inputs tied to V_{CC}, L and G ports in the TRI-STATE mode and tied to ground, all outputs low and tied to ground. The clock monitor and the comparators are disabled.

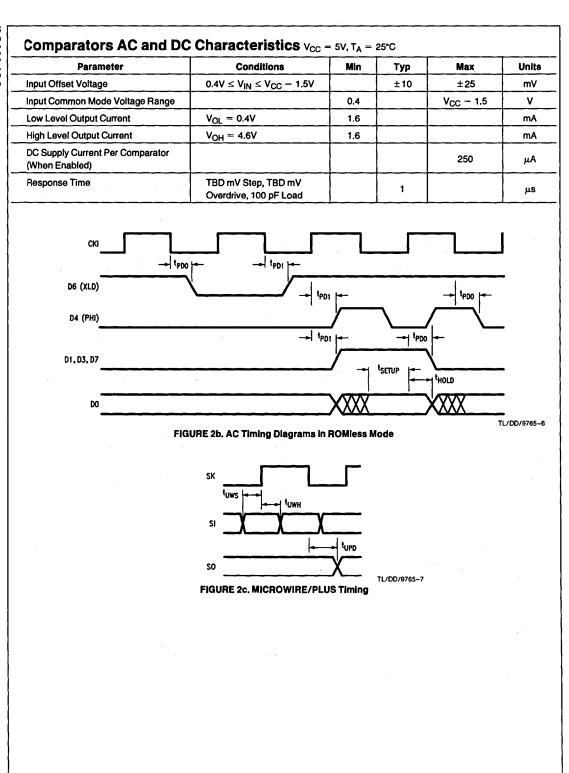
Note 4: Human body model, 100 pF through 1500Ω.

Parameter	Cond	litions	Min	Тур	Ma	ax 🛛	Units
Allowable Sink/Source]		
Current per Pin							
D Outputs (Sink)					1	5	mA
All others					3		mA
Maximum Input Current	T _A = 25°	C			±1	ni I	mA
without Latchup (Note 6)						••• ·	
RAM Retention Voltage, V _r	500 ns Ri		2				v
	and Fall	Time (Min)					
Input Capacitance					7	·	pF
Load Capacitance on D2					10	00	pF
AC Electrical Charact	eristics -	-40°C ≤ T _A ≤	+ 85°C unless c	otherwise spe	cified		
Parameter		Con	ditions	Min	Тур	Max	Unite
Instruction Cycle Time (t _c)							
Crystal, Resonator,		$4V \le V_{CC}$	≤ 6V	1		DC	μs
R/C Oscillator		2.5V ≤ V _{C4}		2.5		DC	μs
		$4V \leq V_{CC}$		3		DC	μs
		2.5V ≤ V _C		7.5		DC	μs
CKI Clock Duty Cycle (Note 5)		f _r = Max		40		60	%
Rise Time (Note 5)		f _r = 10 M⊦	Iz Ext Clock	}		5	ns
Fall Time (Note 5)		f _r = 10 M⊦	Iz Ext Clock			5	ns
Inputs							
^t SETUP		$4V \leq V_{CC}$		200			ns
		2.5V ≤ V _C		500			ns
^t HOLD		$ 4V \le V_{CC}$		60			ns
		$2.5V \leq V_{CO}$	<u>c < 4V</u>	150			ns
Output Propagation Delay		$R_{L} = 2.2k$	$C_{L} = 100 pF$				
^t PD1, ^t PD0							
SO, SK		$4V \leq V_{CC}$				0.7	μs
		$2.5V \leq V_{CO}$	-	1		1.75	μs
All Others		$4V \le V_{CC}$ $2.5V \le V_{CC}$				1 2.5	μs μs
		2.57 3 40		20			+
MICROWIRETM Setup Time (tuws)		1		20 56		}	ns
MICROWIRE Hold Time (t _{UWH}) MICROWIRE Output Propagation D	elav (tupo)			00		220	ns ns
Input Pulse Width							1
Interrupt Input High Time		ł		1)	tc
Interrupt Input Low Time							
Timer Input High Time						l	
Timer Input Low Time		[1		l	
							· ·C

COP888CG

Note 5: Parameter sampled but not 100% tested.

Note 6: Except pin G7: -60 mA to +100 mA (sampled but not 100% tested).



Pin Descriptions

V_{CC} and GND are the power supply pins.

CKI is the clock input. This can come from an R/C generated oscillator, or a crystal oscillator (in conjunction with CKO). See Oscillator Description section.

RESET is the master reset input. See Reset Description section.

The COP888CG contains three bidirectional 8-bit I/O ports (C, G and L), where each individual bit may be independently configured as an input, output or TRI-STATE under program control. Three data memory address locations are allocated for each of these I/O ports. Each I/O port has two associated 8-bit memory mapped registers, the CONFIGU-RATION register and the output DATA register. A memory mapped address is also reserved for the input pins of each I/O port. (See the COP888CG memory map for the various addresses associated with the I/O ports.) *Figure 3* shows the I/O port configurations for the COP888CG. The DATA and CONFIGURATION registers allow for each port bit to be individually configured under software control as shown below:

CONFIGURATION Register	DATA Register	Port Set-Up
0	0	Hi-Z Input
1		(TRI-STATE Output)
0	1	Input with Weak Pull-Up
1	0	Push-Pull Zero Output
1	1	Push-Pull One Output

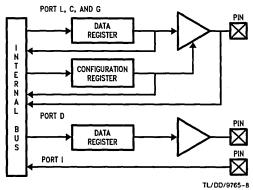


FIGURE 3. I/O Port Configurations

PORT L is an 8-bit I/O port. All L-pins have Schmitt triggers on the inputs.

The Port L supports Multi-Input Wake Up on all eight pins. L1 is used for the UART external clock. L2 and L3 are used for the UART transmit and receive. L4 and L5 are used for the timer input functions T2A and T2B. L6 and L7 are used for the timer input functions T3A and T3B.

The Port L has the following alternate features:

- LO MIWU
- L1 MIWU or CKX
- L2 MIWU or TDX
- L3 MIWU or RDX
- L4 MIWU or T2A

- L5 MIWU or T2B
- L6 MIWU or T3A
- L7 MIWU or T3B

Port G is an 8-bit port with 5 I/O pins (G0, G2–G5), an input pin (G6), and two dedicated output pins (G1 and G7). Pins G0 and G2–G6 all have Schmitt Triggers on their inputs. Pin G1 serves as the dedicated WDOUT WatchDog output, while pin G7 is either input or output depending on the oscillator mask option selected. With the crystal oscillator option selected, G7 serves as the dedicated output pin for the CKO clock output. With the single-pin R/C oscillator mask option selected, G7 serves as a general purpose input pin but is also used to bring the device out of HALT mode with a low to high transition on G7. There are two registers associated with the G Port, a data register and a configuration register. Therefore, each of the 5 I/O bits (G0, G2–G5) can be individually configured under software control.

Since G6 is an input only pin and G7 is the dedicated CKO clock output pin (crystal clock option) or general purpose input (R/C clock option), the associated bits in the data and configuration registers for G6 and G7 are used for special purpose functions as outlined below. Reading the G6 and G7 data bits will return zeros.

Note that the chip will be placed in the HALT mode by writing a "1" to bit 7 of the Port G Data Register. Similarly the chip will be placed in the IDLE mode by writing a "1" to bit 6 of the Port G Data Register.

Writing a "1" to bit 6 of the Port G Configuration Register enables the MICROWIRE/PLUS to operate with the alternate phase of the SK clock. The G7 configuration bit, if set high, enables the clock start up delay after HALT when the R/C clock configuration is used.

	Config Reg.	Data Reg.
G7	CLKDLY	HALT
G6	Alternate SK	IDLE

Port G has the following alternate features:

- G0 INTR (External Interrupt Input)
- G2 T1B (Timer T1 Capture Input)
- G3 T1A (Timer T1 I/O)
- G4 SO (MICROWIRE™ Serial Data Output)
- G5 SK (MICROWIRE Serial Clock)
- G6 SI (MICROWIRE Serial Data Input)

Port G has the following dedicated functions:

- G1 WDOUT WatchDog and/or Clock Monitor dedicated output
- G7 CKO Oscillator dedicated output or general purpose input

PORT I is an eight-bit Hi-Z input port. The 28-pin device does not have a full complement of Port I pins. The unavailable pins are not terminated i.e., they are floating. A read operation for these unterminated pins will return unpredictable values. The user must ensure that the software takes this into account by either masking or restricting the

Pin Descriptions (Continued)

accesses to bit operations. The unterminated Port I pins will draw power only when addressed.

Port I1-I3 are used for Comparator 1. Port I4-I6 are used for Comparator 2.

The Port I has the following alternate features.

- II COMP1-IN (Comparator 1 Negative Input)
- I2 COMP1 + IN (Comparator 1 Positive Input)
- I3 COMP1OUT (Comparator 1 Output)
- 14 COMP2-IN (Comparator 2 Negative Input)
- I5 COMP2+IN (Comparator 2 Positive Input)
- I6 COMP2OUT (Comparator 2 Output)

Port D is an 8-bit output port that is preset high when RESET goes low. The user can tie two or more D port outputs together in order to get a higher drive.

Functional Description

The architecture of the COP888CG is modified Harvard architecture. With the Harvard architecture, the control store program memory (ROM) is separated from the data store memory (RAM). Both ROM and RAM have their own separate addressing space with separate address buses. The COP888CG architecture, though based on Harvard architecture, permits transfer of data from ROM to RAM.

CPU REGISTERS

The CPU can do an 8-bit addition, subtraction, logical or shift operation in one instruction (t_c) cycle time.

There are six CPU registers:

A is the 8-bit Accumulator Register

PC is the 15-bit Program Counter Register

PU is the upper 7 bits of the program counter (PC) PL is the lower 8 bits of the program counter (PC)

B is an 8-bit RAM address pointer, which can be optionally post auto incremented or decremented.

X is an 8-bit alternate RAM address pointer, which can be optionally post auto incremented or decremented.

SP is the 8-bit stack pointer, which points to the subroutine/ interrupt stack (in RAM). The SP is initialized to RAM address 06F with reset.

S is the 8-bit Data Segment Address Register used to extend the lower half of the address range (00 to 7F) into 256 data segments of 128 bytes each.

All the CPU registers are memory mapped with the exception of the Accumulator (A) and the Program Counter (PC).

PROGRAM MEMORY

Program memory for the COP888CG consists of 4096 bytes of ROM. These bytes may hold program instructions or constant data (data tables for the LAID instruction, jump vectors for the JID instruction, and interrupt vectors for the VIS instruction). The program memory is addressed by the 15-bit program counter (PC). All interrupts in the COP888CG vector to program memory location 0FF Hex.

DATA MEMORY

The data memory address space includes the on-chip RAM and data registers, the I/O registers (Configuration, Data and Pin), the control registers, the MICROWIRE/PLUS SIO shift register, and the various registers, and counters associated with the timers (with the exception of the IDLE timer). Data memory is addressed directly by the instruction or indirectly by the B, X, SP pointers and S register.

The COP888CG has 192 bytes of RAM. Sixteen bytes of RAM are mapped as "registers" at addresses 0F0 to 0FF Hex. These registers can be loaded immediately, and also decremented and tested with the DRSZ (decrement register and skip if zero) instruction. The memory pointer registers X, SP, B and S are memory mapped into this space at address locations 0FC to 0FF Hex respectively, with the other registers being available for general usage.

The instruction set of the COP888CG permits any bit in memory to be set, reset or tested. All I/O and registers on the COP888CG (except A and PC) are memory mapped; therefore, I/O bits and register bits can be directly and individually set, reset and tested. The accumulator (A) bits can also be directly and individually tested.

Data Memory Segment RAM Extension

Data memory address 0FF is used as a memory mapped location for the Data Segment Address Register (S) in the COP888CG.

The data store memory is either addressed directly by a single byte address within the instruction, or indirectly relative to the reference of the B, X, or SP pointers (each contains a single-byte address). This single-byte address allows an addressing range of 256 locations from 00 to FF hex. The upper bit of this single-byte address divides the data store memory into two separate sections as outlined previously. With the exception of the RAM register memory from address locations 00F0 to 00FF, all RAM memory is memory mapped with the upper bit of the single-byte address being equal to zero. This allows the upper bit of the single-byte address to determine whether or not the base address range (from 0000 to 00FF) is extended. If this upper bit equals one (representing address range 0080 to 00FF), then address extension does not take place. Alternatively, if this upper bit equals zero, then the data segment extension register S is used to extend the base address range (from 0000 to 007F) from XX00 to XX7F, where XX represents the 8 bits from the S register. Thus the 128-byte data segment extensions are located from addresses 0100 to 017F for data segment 1, 0200 to 027F for data segment 2, etc., up to FF00 to FF7F for data segment 255. The base address range from 0000 to 007F represents data segment 0.

Figure 4 illustrates how the S register data memory extension is used in extending the lower half of the base address range (00 to 7F hex) into 256 data segments of 128 bytes each, with a total addressing range of 32 kbytes from XX00 to XX7F. This organization allows a total of 256 data segments of 128 bytes each with an additional upper base segment of 128 bytes. Furthermore, all addressing modes are available for all data segments. The S register must be changed under program control to move from one data segment (128 bytes) to another. However, the upper base segment (containing the 16 memory registers, I/O registers, control registers, etc.) is always available regardless of the

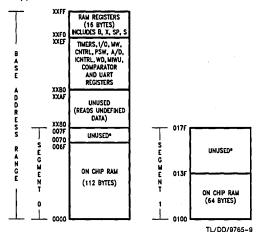
Data Memory Segment RAM Extension (Continued)

contents of the S register, since the upper base segment (address range 0080 to 00FF) is independent of data segment extension.

The instructions that utilize the stack pointer (SP) always reference the stack as part of the base segment (Segment 0), regardless of the contents of the S register. The S register is not changed by these instructions. Consequently, the stack (used with subroutine linkage and interrupts) is always located in the base segment. The stack pointer will be intitialized to point at data memory location 006F as a result of reset.

The 128 bytes of RAM contained in the base segment are split between the lower and upper base segments. The first 116 bytes of RAM are resident from address 0000 to 006F in the lower base segment, while the remaining 16 bytes of RAM represent the 16 data memory registers located at addresses 00F0 to 00FF of the upper base segment. No RAM is located at the upper sixteen addresses (0070 to 007F) of the lower base segment.

Additional RAM beyond these initial 128 bytes, however, will always be memory mapped in groups of 128 bytes (or less) at the data segment address extensions (XX00 to XX7F) of the lower base segment. The additional 64 bytes of RAM in the COP888CG (beyond the initial 128 bytes) are memory mapped at address locations 0100 to 013F hex.



•Reads as all ones.

FIGURE 4. RAM Organization

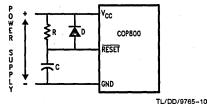
Reset

The RESET input when pulled low initializes the microcontroller. Initialization will occur whenever the RESET input is pulled low. Upon initialization, the data and configuration registers for ports L, G and C are cleared, resulting in these Ports being initialized to the TRI-STATE mode. Pin G1 of the G Port is an exception (as noted below) since pin G1 is

dedicated as the WatchDog and/or Clock Monitor error output pin. Port D is set high. The PC, PSW, ICNTRL, CNTRL, T2CNTRL and T3CNTRL control registers are cleared. The UART registers PSR, ENU (except that TBMT bit is set), ENUR and ENUI are cleared. The Comparator Select Register is cleared. The S register is initialized to zero. The Multi-Input Wakeup registers WKEN, WKEDG and WKPND are cleared. The stack pointer, SP, is initialized to 6F Hex.

The COP888CG comes out of reset with both the Watch-Dog logic and the Clock Monitor detector armed, with the WatchDog service window bits set and the Clock Monitor bit set. The WatchDog and Clock Monitor circuits are inhibited during reset. The WatchDog service window bits being initialized high default to the maximum WatchDog service window of 64k t_C clock cycles. The Clock Monitor bit being initialized high will cause a Clock Monitor error following reset if the clock has not reached the minimum specified frequency at the termination of reset. A Clock Monitor error will cause an active low error output on pic G1. This error output will continue until 16 t_C-32 t_C clock cycles following at which time the G1 output will enter the TRI-STATE mode.

The external RC network shown in *Figure 5* should be used to ensure that the RESET pin is held low until the power supply to the chip stabilizes.



RC > 5 × Power Supply Rise Time FIGURE 5, Recommended Reset Circuit

Oscillator Circuits

The chip can be driven by a clock input on the CKI input pin which can be between DC and 10 MHz. The CKO output clock is on pin G7 (crystal configuration). The CKI input frequency is divided down by 10 to produce the instruction cycle clock $(1/t_c)$.

Figure 6 shows the Crystal and R/C diagrams.

CRYSTAL OSCILLATOR

CKI and CKO can be connected to make a closed loop crystal (or resonator) controlled oscillator.

Table A shows the component values required for various standard crystal values.

R/C OSCILLATOR

By selecting CKI as a single pin oscillator input, a single pin R/C oscillator circuit can be connected to it. CKO is available as a general purpose input, and/or HALT restart input. Table B shows the variation in the oscillator frequencies as functions of the component (R and C) values.

Oscillator Circuits (Continued)

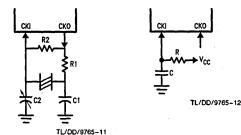


FIGURE 6. Crystal and R/C Oscillator Diagrams

TABLE A. Crystal Oscillator Configuration, $T_A = 25^{\circ}C$

R1 (kΩ)	R2 (ΜΩ)	C1 (pF)	C2 (pF)	CKI Freq (MHz)	Conditions
0	1	30	30-36	10	$V_{CC} = 5V$
0	1	30	30-36	4	$V_{\rm CC} = 2.5 V$
0	1	200	100-150	0.455	$V_{CC} = 2.5V$

TABLE B.	BC Oscillator	Configuration,	T.	=	25°C

R (kΩ)	C (pF)	CKI Freq (MHz)	Instr. Cycle (µs)	Conditions
3.3	82	2.8 to 2.2	3.6 to 4.5	$V_{CC} = 5V$
5.6	100	1.5 to 1.1	6.7 to 9	$V_{CC} = 5V$
6.8	100	1.1 to 0.8	9 to 12.5	$V_{\rm CC} = 2.5V$

Current Drain

The total current drain of the chip depends on:

- 1. Oscillator operation mode-I1
- 2. Internal switching current—I2
- 3. Internal leakage current-I3
- 4. Output source current-I4
- 5. DC current caused by external input not at V_{CC} or GND-15
- 6. Comparator DC supply current when enabled----16

7. Clock Monitor current when enabled-17

Thus the total current drain, It, is given as

To reduce the total current drain, each of the above components must be minimum.

The chip will draw more current as the CKI input frequency increases up to the maximum 10 MHz value. Operating with a crystal network will draw more current than an external square-wave. Switching current, governed by the equation below, can be reduced by lowering voltage and frequency. Leakage current can be reduced by lowering voltage and temperature. The other two items can be reduced by carefully designing the end-user's system.

 $I2 = C \times V \times f$

V = operating voltage

f = CKI frequency

Some sample current drain values at $V_{CC} = 5V$ are:

CKI (MHz)	Inst. Cycle (µs)	It (mA)
10	1	15
3.58	2.8	5.4
2	5	3
0.3	33	0.45
0 (HALT)	— .	0.005

Control Registers

CNTRL Register (Address X'00EE)

The Timer1 (T1) and MICROWIRE/PLUS control register contains the following bits:

	SL1 & SL0	Select the MICROWIRE/PLUS clock divide by $(00 = 2, 01 = 4, 1x = 8)$
	IEDG	External interrupt edge polarity select ($0 = $ Rising edge, $1 = $ Falling edge)
	MSEL	Selects G5 and G4 as MICROWIRE/PLUS signals SK and SO respectively
	T1C0	Timer T1 Start/Stop control in timer modes 1 and 2
		Timer T1 Underflow Interrupt Pending Flag in timer mode 3
	T1C1	Timer T1 mode control bit
	T1C2	Timer T1 mode control bit
_	T1C3	Timer T1 mode control bit
ſ	T1C3 T1C2	TICI TICO MSEL JEDG SLI SLO

T1C3	T1C2	T1C1	T1C0	MSEL	IEDG	SL1	SLO	ł
Bit 7							Bit 0	

PSW Register (Address X'00EF)

The PSW register contains the following select bits:

- GIE Global interrupt enable (enables interrupts)
- EXEN Enable external interrupt
- BUSY MICROWIRE/PLUS busy shifting flag
- EXPND External interrupt pending

T1ENA Timer T1 Interrupt Enable for Timer Underflow or T1A Input capture edge

- T1PNDA Timer T1 Interrupt Pending Flag (Autoreload RA in mode 1, T1 Underflow in Mode 2, T1A capture edge in mode 3)
- C Carry Flag
- HC Half Carry Flag

нс	С	T1PNDA	T1ENA	EXPND	BUSY	EXEN	GIE
Bit 7							Bit 0

The Half-Carry bit is also affected by all the instructions that affect the Carry flag. The SC (Set Carry) and RC (Reset Carry) instructions will respectively set or clear both the carry flags. In addition to the SC and RC instructions, ADC, SUBC, RRC and RLC instructions affect the carry and Half Carry flags.

COP888CG

Control Registers (Continued)

ICNTRL Register (Address X'00E8)

The ICNTRL register contains the following bits:

T1ENB	Timer T1 Interrupt Enable for T1B Input capture
	edge

- T1PNDB Timer T1 Interrupt Pending Flag for T1B capture edge
- µWEN Enable MICROWIRE/PLUS interrupt

µWPND MICROWIRE/PLUS interrupt pending

- TOEN Timer T0 Interrupt Enable (Bit 12 toggle)
- T0PND Timer T0 Interrupt pending
- LPEN L Port Interrupt Enable (Multi-Input Wakeup/Interrupt)

Bit 7 could be used as a flag

Unused	LPEN	TOPND	TOEN	μWPND	μWEN	T1PNDB	T1ENB

Bit 7

T2CNTRL Register (Address X'00C6)

The T2CNTRL register contains the following bits:

- T2ENB Timer T2 Interrupt Enable for T2B Input capture edge
- T2PNDB Timer T2 Interrupt Pending Flag for T2B capture edge
- T2ENA Timer T2 Interrupt Enable for Timer Underflow or T2A Input capture edge

T2PNDA Timer T2 Interrupt Pending Flag (Autoreload RA in mode 1, T2 Underflow in mode 2, T2A capture edge in mode 3)

- T2C0 Timer T2 Start/Stop control in timer modes 1 and 2 Timer T2 Underflow Interrupt Pending Flag in timer mode 3
- T2C1 Timer T2 mode control bit
- T2C2 Timer T2 mode control bit
- T2C3 Timer T2 mode control bit

T2C3	T2C2	T2C1	T2C0	T2PNDA	T2ENA	T2PNDB	T2ENB	
Bit 7							Bit 0	

T3CNTRL Register (Address X'00B6)

The T3CNTRL register contains the following bits:

T3ENB Timer T3 Interrupt Enable for T3B

- T3PNDB Timer T3 Interrupt Pending Flag for T3B pin (T3B capture edge)
- T3ENA Timer T3 Interrupt Enable for Timer Underflow or T3A pin
- T3PNDA Timer T3 Interrupt Pending Flag (Autoload RA in mode 1, T3 Underflow in mode 2, T3a capture edge in mode 3)

ner T3 Start/Stop control in timer modes 1 d 2
mer T3 Underflow Interrupt Pending Flag in ner mode 3
mer T3 mode control bit
mer T3 mode control bit
mer T3 mode control bit

тэсз	T3C2	T3C1	тзсо	T3PNDA	T3ENA	T3PNDB	T3ENB
Bit 7							Bit 0

ROMIess Mode

The COP888CG can address up to 32 kbytes of address space. If at power up, D2 is held at ground, the COP888CG executes from external memory. Port D is used to interface to external program memory. The address comes out in a serial fashion and the data from the external program memory is read back in a serial fashion. The Port D pins perform the following functions.

D0 Shifts in ROM data

Bit 0

- D1 Shifts out lower eight bits of PC
- D2 Places the μ C in the ROMless mode if grounded at reset
- D3 Shifts out upper eight bits of PC
- D4 Data Shift Clock
- D5 HALT Mask Option select pin (D5 = 0) for HALT enable, D5 = 1 for HALT disable)
- D6 Load Clock
- D7 Shifts out recreated Port D data

The most significant bit of the data to come out on the D3 pin is a status signal. It is used by the MOLE development system. This "lost" output port (D0–D7) can be accurately reconstructed with external components as shown in *Figure* δ , providing an accurate emulation.

The 44-pin and 40-pin versions of the COP888CG have a full complement of the D Port pins and can be used in the ROMIess mode.

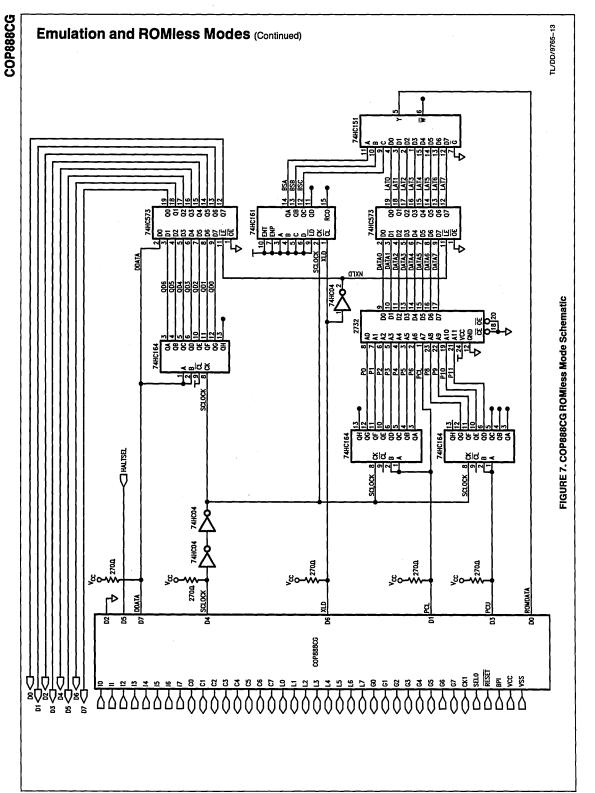
The 28-pin part cannot be used for emulation since it does not have the full complement of 8 D Port pins necessary for entering the ROMIess mode.

Note that in the ROMless mode the D Port is recreated one full CKI clock cycle behind the normal port timings.

Note: Standard parts used in the ROMless mode will operate only at a reduced frequency (to be defined).

The COP888CG device has a spare D pin (D5) in the ROMless mode since only seven pins are required for emulation and recreation. This pin D5 is used in the ROMless mode to enable or disable the HALT mask option feature.

Figure 7 shows the COP888CG ROMless Mode Schematic.



Timers

The COP888CG contains a very versatile set of timers (T0, T1, T2). All timers and associated autoreload/capture registers power up containing random data.

TIMER TO (IDLE TIMER)

The COP888CG supports applications that require maintaining real time and low power with the IDLE mode. This IDLE mode support is furnished by the IDLE timer T0, which is a 16-bit timer. The Timer T0 runs continuously at the fixed rate of the instruction cycle clock, t_c. The user cannot read or write to the IDLE Timer T0, which is a count down timer. The Timer T0 supports the following functions:

Exit out of the Idle Mode (See Idle Mode description) WatchDog logic (See WatchDog description)

Start up delay out of the HALT mode

The IDLE Timer T0 can generate an interrupt when the thirteenth bit toggles. This toggle is latched into the T0PND pending flag, and will occur every 4 ms at the maximum clock frequency ($t_c = 1 \mu_s$). A control flag T0EN allows the interrupt from the thirteenth bit of Timer T0 to be enabled or disabled. Setting T0EN will enable the interrupt, while resetting it will disable the interrupt.

TIMER T1, TIMER T2 AND TIMER T3

The COP888CG has a set of two powerful timer/counter blocks, T1, T2 and T3. The associated features and functioning of a timer block are described by referring to the timer block Tx. Since the three timer blocks, T1, T2 and T3 are identical, all comments are equally applicable to either timer block.

Each timer block consists of a 16-bit timer, Tx, and two supporting 16-bit autoreload/capture registers, RxA and RxB. Each timer block has two pins associated with it, TxA and TxB. The pin TxA supports I/O required by the timer block, while the pin TxB is an input to the timer block. The powerful and flexible timer block allows the COP888CG to easily perform all timer functions with minimal software overhead. The timer block has three operating modes: Processor Independent PWM mode, External Event Counter mode, and Input Capture mode.

The control bits TxC3, TxC2, and TxC1 allow selection of the different modes of operation.

Mode 1. Processor Independent PWM Mode

As the name suggests, this mode allows the COP888CG to generate a PWM signal with very minimal user intervention.

The user only has to define the parameters of the PWM signal (ON time and OFF time). Once begun, the timer block will continuously generate the PWM signal completely independent of the microcontroller. The user software services the timer block only when the PWM parameters require updating.

In this mode the timer Tx counts down at a fixed rate of t_c . Upon every underflow the timer is alternately reloaded with the contents of supporting registers, RxA and RxB. The very first underflow of the timer causes the timer to reload from the register RxA. Subsequent underflows cause the timer to be reloaded from the registers alternately beginning with the register RxB.

The Tx Timer control bits, TxC3, TxC2 and TxC1 set up the timer for PWM mode operation.

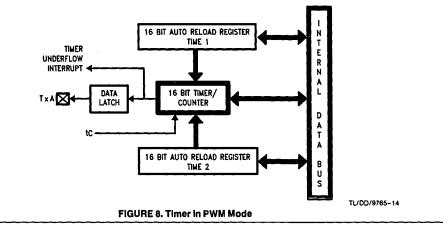
Figure 8 shows a block diagram of the timer in PWM mode. The underflows can be programmed to toggle the TxA output pin. The underflows can also be programmed to generate interrupts.

Underflows from the timer are alternately latched into two pending flags, TxPNDA and TxPNDB. The user must reset these pending flags under software control. Two control enable flags, TxENA and TxENB, allow the interrupts from the timer underflow to be enabled or disabled. Setting the timer enable flag TxENA will cause an interrupt when a timer underflow causes the RxA register to be reloaded into the timer. Setting the timer enable flag TxENB will cause an interrupt when a timer underflow causes the RxB register to be reloaded into the timer. Resetting the timer enable flags will disable the associated interrupts.

Either or both of the timer underflow interrupts may be enabled. This gives the user the flexibility of interrupting once per PWM period on either the rising or falling edge of the PWM output. Alternatively, the user may choose to interrupt on both edges of the PWM output.

Mode 2. External Event Counter Mode

This mode is quite similar to the processor independent PWM mode described above. The main difference is that the timer, Tx, is clocked by the input signal from the TxA pin. The Tx timer control bits, TxC3, TxC2 and TxC1 allow the timer to be clocked either on a positive or negative edge from the TxA pin. Underflows from the timer are latched into the TxPNDA pending flag. Setting the TxENA control flag will cause an interrupt when the timer underflows.



Timers (Continued)

In this mode the input pin TxB can be used as an independent positive edge sensitive interrupt input if the TxENB control flag is set. The occurrence of a positive edge on the TxB input pin is latched into the TxPNDB flag.

Figure 9 shows a block diagram of the timer in External Event Counter mode.

Note: The PWM output is not available in this mode since the TxA pin is being used as the counter input clock.

Mode 3. Input Capture Mode

The COP888CG can precisely measure external frequencies or time external events by placing the timer block, Tx, in the input capture mode.

In this mode, the timer Tx is constantly running at the fixed t_c rate. The two registers, RxA and RxB, act as capture registers. Each register acts in conjunction with a pin. The register RxA acts in conjunction with the TxA pin and the register RxB acts in conjunction with the TxB pin.

The timer value gets copied over into the register when a trigger event occurs on its corresponding pin. Control bits, TxC3, TxC2 and TxC1, allow the trigger events to be specified either as a positive or a negative edge. The trigger condition for each input pin can be specified independently.

The trigger conditions can also be programmed to generate interrupts. The occurrence of the specified trigger condition on the TxA and TxB pins will be respectively latched into the pending flags, TxPNDA and TxPNDB. The control flag TxENA allows the interrupt on TxA to be either enabled or disabled. Setting the TxENA flag enables interrupts to be generated when the selected trigger condition occurs on the TxA pin. Similarly, the flag TxENB controls the interrupts from the TxB pin.

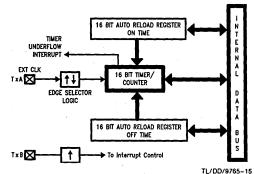


FIGURE 9. Timer in External Event Counter Mode

Underflows from the timer can also be programmed to generate interrupts. Underflows are latched into the timer TxC0 pending flag (the TxC0 control bit serves as the timer underflow interrupt pending flag in the Input Capture mode). Consequently, the TxC0 control bit should be reset when entering the Input Capture mode. The timer underflow interrupt is enabled with the TxENA control flag. When a TxA interrupt occurs in the Input Capture mode, the user must check both the TxPNDA and TxC0 pending flags in order to determine whether a TxA input capture or a timer underflow (or both) caused the interrupt.

Figure 10 shows a block diagram of the timer in Input Capture mode.

TIMER CONTROL FLAGS

The timers T1, T2 and T3 have indentical control structures. The control bits and their functions are summarized below.

- TxC0 Timer Start/Stop control in Modes 1 and 2 (Processor Independent PWM and External Event Counter), where 1 = Start, 0 = Stop Timer Underflow Interrupt Pending Flag in Mode 3 (Input Capture)
- TxPNDA Timer Interrupt Pending Flag
- TxPNDB
 Timer Interrupt Pending Flag

 TxENA
 Timer Interrupt Enable Flag

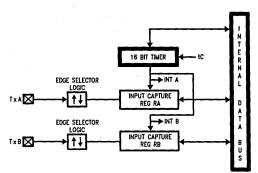
 TxENB
 Timer Interrupt Enable Flag

 1
 Timer Interrupt Enabled

 0
 Timer Interrupt Disabled

	0 = Timer Interrupt Dis
TxC3	Timer mode control
TxC2	Timer mode control

TxC1 Timer mode control



TL/DD/9765-16

FIGURE 10. Timer in Input Capture Mode

COP888CG

TxC3	TxC2	TxC1	Timer Mode	Interrupt A Source	Interrupt B Source	Timer Counts On
0	0	0	MODE 2 (External Event Counter)	Timer Underflow	Pos. TxB Edge	TxA Pos. Edge
0	0	1	MODE 2 (External Event Counter)	Timer Underflow	Pos. TxB Edge	TxA Neg. Edge
1	0	1	MODE 1 (PWM) TxA Toggle	Autoreload RA	Autoreload RB	tc
1	0	0	MODE 1 (PWM) No TxA Toggle	Autoreload RA	Autoreload RB	tc
0	1	0	MODE 3 (Capture) Captures: TxA Pos. Edge TxB Pos. Edge	Pos. TxA Edge or Timer Underflow	Pos. TxB Edge	tc
1	1	0	MODE 3 (Capture) Captures: TxA Pos. Edge TxB Neg. Edge	Pos. TxA Edge or Timer Underflow	Neg. TxB Edge	tc
0	1	1	MODE 3 (Capture) Captures: TxA Neg. Edge TxB Pos. Edge	Neg. TxB Edge or Timer Underflow	Pos. TxB Edge	t _c
1	1	1	MODE 3 (Capture) Captures: TxA Neg. Edge TxB Neg. Edge	Neg. TxA Edge or Timer Underflow	Neg. TxB Edge	t _c

Power Save Modes

The COP888CG offers the user two power save modes of operation: HALT and IDLE. In the HALT mode, all microcontroller activities are stopped. In the IDLE mode, the onboard oscillator circuitry the WatchDog logic, the Clock Monitor and timer T0 are active but all other microcontroller activities are stopped. In either mode, all on-board RAM, registers, I/O states, and timers (with the exception of T0) are unaltered.

HALT MODE

The COP888CG is placed in the HALT mode by writing a "1" to the HALT flag (G7 data bit). All microcontroller activities, including the clock and timers, are stopped. The WatchDog logic on the COP888CG is disabled during the HALT mode. However, the clock monitor circuitry if enabled remains active. In the HALT mode, the power requirements of the COP888CG are minimal and the applied voltage (V_{CC}) may be decreased to V_r (V_r = 2.0V) without altering the state of the machine.

The COP888CG supports three different ways of exiting the HALT mode. The first method of exiting the HALT mode is with the Multi-Input Wakeup feature on the L port. The second method is with a low to high transition on the CKO (G7) pin. This method precludes the use of the crystal clock configuration (since CKO becomes a dedicated output), and so may be used with an RC clock configuration. The third method of exiting the HALT mode is by pulling the RESET pin low.

Since a crystal or ceramic resonator may be selected as the oscillator, the Wakeup signal is not allowed to start the chip running immediately since crystal oscillators and ceramic resonators have a delayed start up time to reach full amplitude and frequency stability. The IDLE timer is used to generate a fixed delay to ensure that the oscillator has indeed stabilized before allowing instruction execution. In this case, upon detecting a valid Wakeup signal, only the oscillator circuitry is enabled. The IDLE timer is loaded with a value of 256 and is clocked with the t_c instruction cycle clock. The t_c clock is derived by dividing the oscillator clock down by a factor of 10. The Schmitt trigger following the CKI inverter on the chip ensures that the IDLE timer is clocked only when the oscillator has a sufficiently large amplitude to meet the Schmitt trigger specifications. This Schmitt trigger is not part of the oscillator closed loop. The startup timeout from the IDLE timer enables the clock signals to be routed to the rest of the chip.

If an RC clock option is being used, the fixed delay is introduced optionally. A control bit, CLKDLY, mapped as configuration bit G7, controls whether the delay is to be introduced or not. The delay is included if CLKDLY is set, and excluded if CLKDLY is reset. The CLKDLY bit is cleared on reset.

The COP888CG has two mask options associated with the HALT mode. The first mask option enables the HALT mode

Power Save Modes (Continued)

feature, while the second mask option disables the HALT mode. With the HALT mode enable mask option, the COP888CG will enter and exit the HALT mode as described above. With the HALT disable mask option, the COP888CG cannot be placed in the HALT mode (writing a "1" to the HALT flag will have no effect).

The WatchDog detector circuit is inhibited during the HALT mode. However, the clock monitor circuit if enabled remains active during HALT mode in order to ensure a clock monitor error if the COP888CG inadvertently enters the HALT mode as a result of a runaway program or power glitch.

IDLE MODE

The COP888CG is placed in the IDLE mode by writing a "1" to the IDLE flag (G6 data bit). In this mode, all activities, except the associated on-board oscillator circuitry, the WatchDog logic, the clock monitor and the IDLE Timer T0, are stopped. The power supply requirements of the micro-controller in this mode of operation are typically around 30% of normal power requirement of the microcontroller.

As with the HALT mode, the COP888CG can be returned to normal operation with a reset, or with a Multi-Input Wakeup from the L Port. Alternately, the microcontroller resumes normal operation from the IDLE mode when the thirteenth bit (representing 4.096 ms at internal clock frequency of 1 MHz, $t_c = 1 \mu_S$) of the IDLE Timer toggles.

This toggle condition of the thirteenth bit of the IDLE Timer T0 is latched into the T0PND pending flag.

The user has the option of being interrupted with a transition on the thirteenth bit of the IDLE Timer T0. The interrupt can be enabled or disabled via the T0EN control bit. Setting the T0EN flag enables the interrupt and vice versa.

The user can enter the IDLE mode with the Timer T0 interrupt enabled. In this case, when the T0PND bit gets set, the COP888CG will first execute the Timer T0 interrupt service routine and then return to the instruction following the "Enter Idle Mode" instruction.

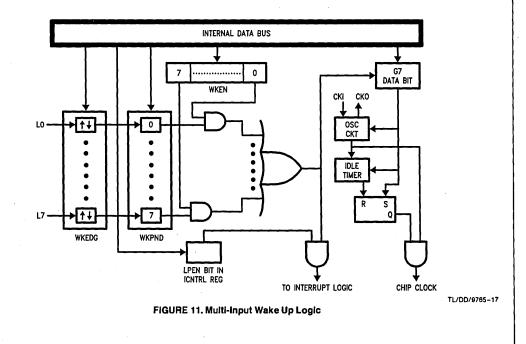
Alternatively, the user can enter the IDLE mode with the IDLE Timer T0 interrupt disabled. In this case, the COP868CG will resume normal operation with the instruction immediately following the "Enter IDLE Mode" instruction.

Note: It is necessary to program two NOP instructions following both the set HALT mode and set IDLE mode instructions. These NOP instructions are necessary to allow clock resynchronization following the HALT or IDLE modes.

Multi-Input Wakeup

The Multi-Input Wakeup feature is ued to return (wakeup) the COP888CG from either the HALT or IDLE modes. Alternately Multi-Input Wakeup/Interrupt feature may also be used to generate up to 8 edge selectable external interrupts.

Figure 11 shows the Multi-Input Wakeup logic for the COP888CG microcontroller.



Multi-Input Wakeup (Continued)

The Multi-Input Wakeup feature utilizes the L Port. The user selects which particular L port bit (or combination of L Port bits) will cause the COP888CG to exit the HALT or IDLE modes. The selection is done through the Reg: WKEN. The Reg: WKEN is an 8-bit read/write register, which contains a control bit for every L port bit. Setting a particular WKEN bit enables a Wakeup from the associated L port pin.

The user can select whether the trigger condition on the selected L Port pin is going to be either a positive edge (low to high transition) or a negative edge (high to low transition). This selection is made via the Reg: WKEDG, which is an 8-bit control register with a bit assigned to each L Port pin. Setting the control bit will select the trigger condition to be a negative edge on that particular L Port pin. Resetting the bit selects the trigger condition to be a positive edge. Changing an edge select entails several steps in order to avoid a pseudo Wakeup condition as a result of the edge change. First, the associated WKEN bit should be reset, followed by the edge select change in WKEDG. Next, the associated WKEN bit should be cleared, followed by the associated WKEN bit being re-enabled.

An example may serve to clarify this procedure. Suppose we wish to change the edge select from positive (low going high) to negative (high going low) for L Port bit 5, where bit 5 has previously been enabled for an input interrupt. The program would be as follows:

RBIT	5,	WKEN
SBIT	5,	WKEDG
RBIT	5,	WKPND
SBIT	5,	WKEN
SBIT	5,	WKEN

If the L port bits have been used as outputs and then changed to inputs with Multi-Input Wakeup/Interrupt, a safety procedure should also be followed to avoid inherited pseudo wakeup conditions. After the selected L port bits have been changed from output to input but before the associated WKEN bits are enabled, the associated edge select bits in WKEDG should be set or reset for the desired edge selects, followed by the associated WKPND bits being cleared.

This same procedure should be used following reset, since the L port inputs are left floating as a result of reset.

The occurrence of the selected trigger condition for Multi-Input Wakeup is latched into a pending register called WKPND. The respective bits of the WKPND register will be set on the occurrence of the selected trigger edge on the corresponding Port L pin. The user has the responsibility of clearing these pending flags. Since WKPND is a pending register for the occurrence of selected wakeup conditions, the COP888CG will not enter the HALT mode if any Wakeup bit is both enabled and pending. Consequently, the user has the responsibility of clearing the pending flags before attempting to enter the HALT mode.

WKEN, WKPND and WKEDG are all read/write registers, and are cleared at reset.

PORT L INTERRUPTS

Port L provides the user with an additional eight fully selectable, edge sensitive interrupts which are all vectored into the same service subroutine.

The interrupt from Port L shares logic with the wake up circuitry. The register WKEN allows interrupts from Port L to be individually enabled or disabled. The register WKEDG specifies the trigger condition to be either a positive or a negative edge. Finally, the register WKPND latches in the pending trigger conditions.

The GIE (Global Interrupt Enable) bit enables the interrupt function.

A control flag, LPEN, functions as a global interrupt enable for Port L interrupts. Setting the LPEN flag will enable interrupts and vice versa. A separate global pending flag is not needed since the register WKPND is adequate.

Since Port L is also used for waking the COP888CG out of the HALT or IDLE modes, the user can elect to exit the HALT or IDLE modes either with or without the interrupt enabled. If he elects to disable the interrupt, then the COP888CG will restart execution from the instruction immediately following the instruction that placed the microcontroller in the HALT or IDLE modes. In the other case, the COP888CG will first execute the interrupt service routine and then revert to normal operation.

The Wakeup signal will not start the chip running immediately since crystal oscillators or ceramic resonators have a finite start up time. The IDLE Timer (T0) generates a fixed delay to ensure that the oscillator has indeed stabilized before allowing the COP888CG to execute instructions. In this case, upon detecting a valid Wakeup signal, only the oscillator circuitry and the IDLE Timer T0 are enabled. The IDLE Timer is loaded with a value of 256 and is clocked from the t_c instruction cycle clock. The t_c clock is derived by dividing down the oscillator clock by a factor of 10. A Schmitt trigger following the CKI on-chip inverter ensures that the IDLE timer is clocked only when the oscillator has a sufficiently large amplitude to meet the Schmitt trigger specifications. This Schmitt trigger is not part of the oscillator closed loop. The startup timeout from the IDLE timer enables the clock signals to be routed to the rest of the chip.

COP888CG

Multi-Input Wakeup (Continued)

If the RC clock option is used, the fixed delay is under software control. A control flag, CLKDLY, in the G7 configuration bit allows the clock start up delay to be optionally inserted. Setting CLKDLY flag high will cause clock start up delay to be inserted and resetting it will exclude the clock start up delay. The CLKDLY flag is cleared during reset, so the clock start up delay is not present following reset with the RC clock options.

UART

The COP888CG contains a full-duplex software programmable UART. The UART (*Figure 12*) consists of a transmit shift register, a receiver shift register and seven addressable registers, as follows: a transmit buffer register (TBUF), a receiver buffer register (RBUF), a UART control and status register (ENU), a UART receive control and status register (ENUR), a UART interrupt and clock source register (ENUI), a prescaler select register (PSR) and baud (BAUD) register. The ENU register contains flags for transmit and receive functions; this register also determines the length of the data frame (7, 8 or 9 bits), the value of the ninth bit in transmission, and parity selection bits. The ENUR register flags framming, data overrun and parity errors while the UART is receiving.

Other functions of the ENUR register include saving the ninth bit received in the data frame, enabling or disabling the UART's attention mode of operation and providing additional receiver/transmitter status information via RCVG and XMTG bits. The determination of an internal or external clock source is done by the ENUI register, as well as selecting the number of stop bits and enabling or disabling transmit and receive interrupts. A control flag in this register can also select the UART mode of operation: asynchronous or synchronous.

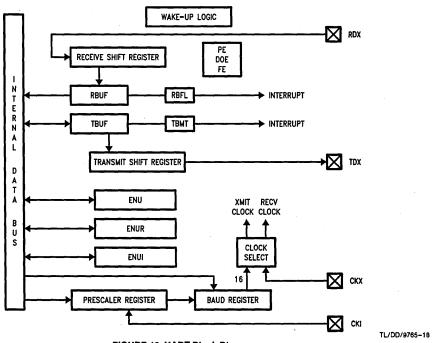


FIGURE 12. UART Block Diagram

UART (Continued)

UART CONTROL AND STATUS REGISTERS

The operation of the UART is programmed through three registers: ENU, ENUR and ENUI. The function of the individual bits in these registers is as follows:

ENU-UART Control and Status Register (Address at OBA)

Į	PEN	PSEL1	XBIT9/	CHL1	CHLO	ERR	RBFL	твмт
			PSEL0					
ł	ORW	ORW	0RW	0RW	ORW	0R	0R	1R

Bit 7

ENUR-UART Receive Control and Status Register (Address at 0BB)

DOE 0RD	FE 0RD	PE 0RD	SPARE 0RW*	1	RCVG 0R
Bit7					Bit0

ENUI-UART Interrupt and Clock Source Register (Address at 0BC)

		XRCLK ORW		ETI ORW
Bit7	 			Bit0

Bit7

*Bit is not used.

Bit is cleared on reset. 0

Bit is set to one on reset.

R Bit is read-only; it cannot be written by software.

- **BW** Bit is read/write.
- Bit is cleared on read; when read by software as a one, it is cleared D automatically. Writing to the bit does not affect its state.

DESCRIPTION OF UART REGISTER BITS

ENU-UART CONTROL AND STATUS REGISTER

TBMT: This bit is set when the UART transfers a byte of data from the TBUF register into the TSFT register for transmission. It is automatically reset when software writes into the TBUF register.

RBFL: This bit is set when the UART has received a complete character and has copied it into the RBUF register. It is automatically reset when software reads the character from RBUF.

ERR: This bit is a global UART error flag which gets set if any or a combination of the errors (DOE, FE, PE) occur.

CHL1, CHL0: These bits select the character frame format. Parity is not included and is generated/verified by hardware.

CHL1 = 0, CHL0 = 0	The frame contains eight data bits.
CHL1 = 0, CHL0 = 1	The frame contains seven data
	bits.

CHL1 = 1, CHL0 = 0CHL1 = 1, CHL0 = 1 The frame contains nine data bits. Loopback Mode selected. Transmitter output internally looped back to receiver input. Nine bit framing format is used.

XBIT9/PSEL0: Programs the ninth bit for transmission when the UART is operating with nine data bits per frame. For seven or eight data bits per frame, this bit in conjunction with PSEL1 selects parity.

PSEL1, PSEL0: Parity select bits.

PSEL1 = 0, PSEL0 = 0 Odd Parity (if Parity enabled) PSEL1 = 0, PSEL0 = 1 Even Parity (if Parity enabled) PSEL1 = 1, PSEL0 = 0 Mark(1) (if Parity enabled) PSEL1 = 1, PSEL0 = 1Space(0) (if Parity enabled) PEN: This bit enables/disables Parity (7- and 8-bit modes

only). PEN = 0 Parity disabled.

PEN = 1 Parity enabled.

Bit 0

ENUR-UART RECEIVE CONTROL AND STATUS REGISTER

RCVG: This bit is set high whenever a framing error occurs and goes low when RDX goes high.

XMTG: This bit is set to indicate that the UART is transmitting. It gets reset at the end of the last frame (end of last Stop bit).

ATTN: ATTENTION Mode is enabled while this bit is set. This bit is cleared automatically on receiving a character with data bit nine set.

RBIT9: Contains the ninth data bit received when the UART is operating with nine data bits per frame.

SPARE: Reserved for future use.

- PE: Flags a Parity Error.
- PE = 0 Indicates no Parity Error has been detected since the last time the ENUR register was read.
- PE = 1 Indicates the occurence of a Parity Error.
- FE: Flags a Framing Error.
- FE = 0 Indicates no Framing Error has been detected since the last time the ENUR register was read.
- FE = 1 Indicates the occurence of a Framing Error.
- DOE: Flags a Data Overrun Error.
- DOE = 0Indicates no Data Overrun Error has been detected since the last time the ENUR register was read
- DOE = 1Indicates the occurence of a Data Overrun Error.

ENUI-UART INTERRUPT AND CLOCK SOURCE REGISTER

ETI: This bit enables/disables interrupt from the transmitter section.

- Interrupt from the transmitter is disabled. FTI = 0
- Interrupt from the transmitter is enabled. ETI = 1
- ERI: This bit enables/disables interrupt from the receiver section.
- ERI = 0Interrupt from the receiver is disabled.
- Interrupt from the receiver is enabled. ERI = 1

XTCLK: This bit selects the clock source for the transmittersection.

- XTCLK = 0The clock source is selected through the PSR and BAUD registers.
- XTCLK = 1 Signal on CKX (L1) pin is used as the clock.
- XRCLK: This bit selects the clock source for the receiver section.
- XRCLK = 0The clock source is selected through the PSR and BAUD registers.

Signal on CKX (L1) pin is used as the clock. XRCLK = 1

- SSEL: UART mode select.
- SSEL = 0 Asynchronous Mode.
- SSEL = 1 Synchronous Mode.

UART (Continued)

ETDX: TDX (UART Transmit Pin) is the alternate function assigned to Port L pin L2; it is selected by setting ETDX bit. To simulate line break generation, software should reset ETDX bit and output logic zero to TDX pin through Port L data and configuration registers.

STP78: This bit is set to program the last Stop bit to be 7/8th of a bit in length.

STP2: This bit programs the number of Stop bits to be transmitted.

STP2 = 0 One Stop bit transmitted. STP2 = 1 Two Stop bits transmitted.

Associated I/O Pins

Data is transmitted on the TDX pin and received on the RDX pin. TDX is the alternate function assigned to Port L pin L2; it is selected by setting ETDX (in the ENUI register) to one. RDX is an inherent function of Port L pin L3, requiring no setup.

The baud rate clock for the UART can be generated onchip, or can be taken from an external source. Port L pin L1 (CKX) is the external clock I/O pin. The CKX pin can be either an input or an output, as determined by Port L Configuration and Data registers (Bit 1). As an input, it accepts a clock signal which may be selected to drive the transmitter and/or receiver. As an output, it presents the internal Baud Rate Generator output.

UART Operation

The UART has two modes of operation: asynchronous mode and synchronous mode.

ASYNCHRONOUS MODE

This mode is selected by resetting the SSEL (in the ENUI register) bit to zero. The input frequency to the UART is 16 times the baud rate.

The TSFT and TBUF registers double-buffer data for transmission. While TSFT is shifting out the current character on the TDX pin, the TBUF register may be loaded by software with the next byte to be transmitted. When TSFT finishes transmitting the current character the contents of TBUF are transferred to the TSFT register and the Transmit Buffer Empty Flag (TBMT in the ENU register) is set. The TBMT flag is automatically reset by the UART when software loads a new character into the TBUF register. There is also the XMTG bit which is set to indicate that the UART is transmitting. This bit gets reset at the end of the last frame (end of last Stop bit). TBUF is a read/write register.

The RSFT and RBUF registers double-buffer data being received. The UART receiver continually monitors the signal on the RDX pin for a low level to detect the beginning of a Start bit. Upon sensing this low level, it waits for half a bit time and samples again. If the RDX pin is still low, the receiver considers this to be a valid Start bit, and the remaining bits in the character frame are each sampled a single time, at the mid-bit position. Serial data input on the RDX pin is shifted into the RSFT register. Upon receiving the complete character, the contents of the RSFT register are copied into the RBUF register and the Received Buffer Full Flag (RBFL) is set. RBFL is automatically reset when software reads the character from the RBUF register. RBUF is a read only register. There is also the RCVG bit which is set high when a framing error occurs and goes low once RDX goes high. TBMT, XMTG, RBFL and RCVG are read only bits.

SYNCHRONOUS MODE

In this mode data is transferred synchronously with the clock. Data is transmitted on the rising edge and received on the falling edge of the synchronous clock.

This mode is selected by setting SSEL bit in the ENUI register. The input frequency to the UART is the same as the baud rate.

When an external clock input is selected at the CKX pin, data transmit and receive are performed synchronously with this clock through TDX/RDX pins.

If data transmit and receive are selected with the CKX pin as clock output, the μ C generates the synchronous clock output at the CKX pin. The internal baud rate generator is used to produce the synchronous clock. Data transmit and receive are performed synchronously with this clock.

FRAMING FORMATS

The UART supports several serial framing formats (*Figure 13*). The format is selected using control bits in the ENU, ENUR and ENUI registers.

The first format (1, 1a, 1b, 1c) for data transmission (CHL0 = 1, CHL1 = 0) consists of Start bit, seven Data bits (excluding parity) and 7/8, one or two Stop bits. In applications using parity, the parity bit is generated and verified by hardware.

The second format (CHL0 = 0, CHL1 = 0) consists of one Start bit, eight Data bits (excluding parity) and 7/8, one or two Stop bits. Parity bit is generated and verified by hardware.

The third format for transmission (CHL0 = 0, CHL1 = 1) consists of one Start bit, nine Data bits and 7/8, one or two Stop bits. This format also supports the UART "ATTEN-TION" feature. When operating in this format, all eight bits of TBUF and RBUF are used for data. The ninth data bit is transmitted and received using two bits in the ENU and ENUR registers, called XBIT9 and RBIT9. RBIT9 is a read only bit. Parity is not generated or verified in this mode.

For any of the above framing formats, the last Stop bit can be programmed to be 7/8th of a bit in length. If two Stop bits are selected and the 7/8th bit is set (selected), the second Stop bit will be 7/8th of a bit in length.

The parity is enabled/disabled by PEN bit located in the ENU register. Parity is selected for 7- and 8-bit modes only. If parity is enabled (PEN = 1), the parity selection is then performed by PSEL0 and PSEL1 bits located in the ENU register.

Note that the XBIT9/PSEL0 bit located in the ENU register serves two mutually exclusive functions. This bit programs the ninth bit for transmission when the UART is operating with nine data bits per frame. There is no parity selection in this framing format. For other framing formats XBIT9 is not needed and the bit is PSEL0 used in conjunction with PSEL1 to select parity.

The frame formats for the receiver differ from the transmitter in the number of Stop bits required. The receiver only requires one Stop bit in a frame, regardless of the setting of the Stop bit selection bits in the control register. Note that an implicit assumption is made for full duplex UART operation that the framing formats are the same for the transmitter and receiver.

COP888CG

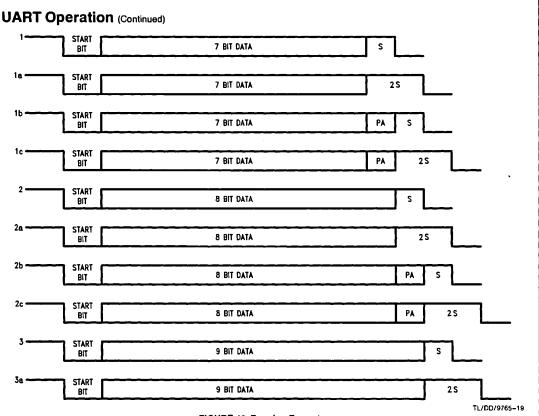


FIGURE 13. Framing Formats

UART INTERRUPTS

The UART is capable of generating interrupts. Interrupts are generated on Receive Buffer Full and Transmit Buffer Empty. Both interrupts have individual interrupt vectors. Two bytes of program memory space are reserved for each interrupt vector. The two vectors are located at addresses 0xEC to 0xEF Hex in the program memory space. The interrupts can be individually enabled or disabled using Enable Transmit Interrupt (ETI) and Enable Receive Interrupt (ERI) bits in the ENUI register.

The interrupt from the Transmitter is set pending, and remains pending, as long as both the TBMT and ETI bits are set. To remove this interrupt, software must either clear the ETI bit or write to the TBUF register (thus clearing the TBMT bit).

The interrupt from the receiver is set pending, and remains pending, as long as both the RBFL and ERI bits are set. To remove this interrupt, software must either clear the ERI bit or read from the RBUF register (thus clearing the RBFL bit).

Baud Clock Generation

The clock inputs to the transmitter and receiver sections of the UART can be individually selected to come either from an external source at the CKX pin (port L, pin L1) or from a source selected in the PSR and BAUD registers. Internally, the basic baud clock is created from the oscillator frequency through a two-stage divider chain consisting of a 1–16 (increments of 0.5) prescaler and an 11-bit binary counter. (*Figure 14*) The divide factors are specified through two read/write registers shown in *Figure 15*. Note that the 11-bit Baud Rate Divisor spills over into the Prescaler Select Register (PSR). PSR is cleared upon reset.

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As shown in Table I, a Prescaler Factor of 0 corresponds to NO CLOCK. NO CLOCK condition is the UART power down mode where the UART clock is turned off for power saving purpose. The user must also turn the UART clock off when a different baud rate is chosen.

The correspondences between the 5-bit Prescaler Select and Prescaler factors are shown in Table I. Therer are many ways to calculate the two divisor factors, but one particularly effective method would be to achieve a 1.8432 MHz prescaler output is then used to drive the software programmable baud rate counter to create a x16 clock for the following baud rates: 110, 134.5, 150, 300, 600, 1200, 1800, 2400, 3600, 4800, 7200, 9600, 19200 and 38400 (Table II). Other baud rates may be created by using appropriate divisors. The x16 clock is then divided by 16 to provide the rate for the serial shift registers of the transmitter and receiver.

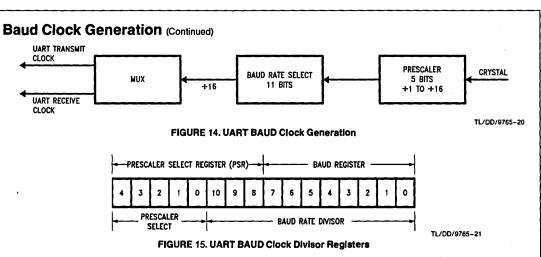


TABLE I. Prescaler Factors

Prescaler	Prescaler
Select	Factor
00000	NO CLOCK
00001	1
00010	1.5
00011	2
00100	2.5
00101	3
00110	3.5
00111	4
01000	4.5
01001	5
01010	5.5
01011	6
01100	6.5
01101	7
01110	7.5
01111	8
10000	8.5
10001	9
10010	9.5
10011	10
10100	10.5
10101	11
10110	11.5
10111	12
11000	12.5
11001	13
11010	13.5
11011	14
11100	14.5
11101	15
11110	15.5
11111	16

TABLE II. Baud Rate Divisors (1.8432 MHz Prescaler Output)

(110 +02 1111=1	restater suputy
Baud Rate	Baud Rate Divisor — 1 (N-1)
110 (110.03)	1046
134.5 (134.58)	855
150	767
300	383
600	191
1200	95
1800	63
2400	47
3600	31
4800	23
7200	15
9600	11
19200	5
38400	2

The entries in Table II assume a prescaler output of 1.8432 MHz. In the asynchronous mode the baud rate could be as high as 625k.

As an example, considering the Asynchronous Mode and a CKI clock of 4.608 MHz, the prescaler factor selected is:

4.608/1.8432 = 2.5

The 2.5 entry is available in Table I. The 1.8432 MHz prescaler output is then used with proper Baud Rate Divisor (Table II) to obtain different baud rates. For a baud rate of 19200 e.g., the entry in Table II is 5.

N - 1 = 5 (N - 1 is the value from Table II)

N = 6 (N is the Baud Rate Divisor)

Baud Rate = $1.8432 \text{ MHz}/(16 \times 6) = 19200$

The divide by 16 is performed because in the asynchronous mode, the input frequency to the UART is 16 times the baud rate. The equation to calculate baud rates is given below. The actual Baud Rate may be found from:

 $BR = Fc/(16 \times N \times P)$

Baud Clock Generation (Continued)

Where:

BR is the Baud Rate

Fc is the CKI frequency

N is the Baud Rate Divisor (Table II).

P is the Prescaler Divide Factor selected by the value in the Prescaler Select Register (Table I)

Note: In the Synchronous Mode, the divisor 16 is replaced by one.

Example:

Asynchronous Mode:

Crystal Frequency = 5 MHz

Desired baud rate = 9600

Using the above equation N \times P can be calculated first.

 $N \times P = (5 \times 10^6)/(16 \times 9600) = 32.552$

Now 32.552 is divided by each Prescaler Factor (Table II) to obtain a value closest to an integer. This factor happens to be 6.5 (P = 6.5).

N = 32.552/6.5 = 5.008 (N = 5)

The programmed value (from Table II) should be 4 (N - 1). Using the above values calculated for N and P:

 $BR = (5 \times 10^6)/(16 \times 5 \times 6.5) = 9615.384$ % error = (9615.385 - 9600)/9600 = 0.16

Effect of HALT/IDLE

The UART logic is reinitialized when either the HALT or IDLE modes are entered. This reinitialization sets the TBMT flag and resets all read only bits in the UART control and status registers. Read/Write bits remain unchanged. The Transmit Buffer (TBUF) is not affected, but the Transmit Shift register (TSFT) bits are set to one. The receiver registers RBUF and RSFT are not affected.

The μ C will exit from the HALT/IDLE modes when the Start bit of a character is detected at the RDX (L3) pin. This feature is obtained by using the Multi-Input Wakeup scheme provided on the μ C.

Before entering the HALT or IDLE modes the user program must select the Wakeup source to be on the RDX pin. This selection is done by setting bit 3 of WKEN (Wakeup Enable) register. The Wakeup trigger condition is then selected to be high to low transition. This is done via the WKEDG register (Bit 3 is zero.)

If the microcontroller is halted and crystal oscillator is used, the Wakeup signal will not start the chip running immediately because of the finite start up time requirement of the crystal oscillator. The idle timer (T0) generates a fixed delay to ensure that the oscillator has indeed stabilized before allowing the μC to execute code. The user has to consider this delay when data transfer is expected immediately after exiting the HALT mode.

Diagnostic

Bits CHARL0 and CHARL1 in the ENU register provide a loopback feature for diagnostic testing of the UART. When these bits are set to one, the following occur: The receiver input pin (RDX) is internally connected to the transmitter output pin (TDX); the output of the Transmitter Shift Register is "looped back" into the Receive Shift Register input. In this mode, data that is transmitted is immediately received. This feature allows the processor to verify the transmit and receive data paths of the UART.

Note that the framing format for this mode is the nine bit format; one Start bit, nine data bits, and 7/8, one or two Stop bits. Parity is not generated or verified in this mode.

Attention Mode

The UART Receiver section supports an alternate mode of operation, referred to as ATTENTION Mode. This mode of operation is selected by the ATTN bit in the ENUR register. The data format for transmission must also be selected as having nine Data bits and either 7/8, one or two Stop bits.

The ATTENTION mode of operation is intended for use in networking the COP888CG with other processors. Typically in such environments the messages consists of device addresses, indicating which of several destinations should receive them, and the actual data. This Mode supports a scheme in which addresses are flagged by having the ninth bit of the data field set to a 1. If the ninth bit is reset to a zero the byte is a Data byte.

While in ATTENTION mode, the UART monitors the communication flow, but ignores all characters until an address character is received. Upon receiving an address character, the UART signals that the character is ready by setting the RBFL flag, which in turn interrupts the processor if UART Receiver interrupts are enabled. The ATTN bit is also cleared automatically at this point, so that data characters as well as address characters are recognized. Software examines the contents of the RBUF and responds by deciding either to accept the subsequent data stream (by leaving the ATTN bit reset) or to wait until the next address character is seen (by setting the ATTN bit again).

Operation of the UART Transmitter is not affected by selection of this Mode. The value of the ninth bit to be transmitted is programmed by setting XBIT9 appropriately. The value of the ninth bit received is obtained by reading RBIT9. Since this bit is located in ENUR register where the error flags reside, a bit operation on it will reset the error flags.

Comparators

The COP888CG contains two differential comparators, each with a pair of inputs (positive and negative) and an output. Ports 11–13 and 14–16 are used for the comparators. The following is the Port I assignment:

- 11 Comparator1 negative input
- 12 Comparator1 positive input
- 13 Comparator1 output
- 14 Comparator2 negative input
- 15 Comparator2 positive input
- 16 Comparator2 output

A Comparator Select Register (CMPSL) is used to enable the comparators, read the outputs of the comparators internally, and enable the outputs of the comparators to the pins. Two control bits (enable and output enable) and one result bit are associated with each comparator. The comparator result bits (CMP1RD and CMP2RD) are read only bits which will read as zero if the associated comparator is not enabled. The Comparator Select Register is cleared with reset, resulting in the comparators being disabled. The comparators should also be disabled before entering either the HALT or IDLE modes in order to save power. The configuration of the CMPSL register is as follows:

Comparators (Continued)

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CMPSL REGISTER (ADDRESS X'00B7)

The CMPSL register contains the following bits:

- CMP1EN Enable comparator 1
- CMP1RD Comparator 1 result (this is a read only bit, which will read as 0 if the comparator is not enabled)
- CMP10E Selects pin I3 as comparator 1 output provided that CMPIEN is set to enable the comparator
- CMP2EN Enable comparator 2
- CMP2RD Comparator 2 result (this is a read only bit, which will read as 0 if the comparator is not enabled)
- CMP20E Selects pin I6 as comparator 2 output provided that CMP2EN is set to enable the comparator

Unused CMP20E CMP2RD CMP2EN CMP10E CMP1RD CMP1EN Unused Bit 7 Bit 0

Note that the two unused bits of CMPSL may be used as software flags.

Comparator outputs have the same spec as Ports L and G except that the rise and fall times are symmetrical.

Interrupts

The COP888CG supports a vectored interrupt scheme. It supports a total of fourteen interrupt sources. The following table lists all the possible COP888CG interrupt sources, their arbitration ranking and the memory locations reserved for the interrupt vector for each source.

Two bytes of program memory space are reserved for each interrupt source. All interrupt sources except the software interrupt are maskable. Each of the maskable interrupts have an Enable bit and a Pending bit. A maskable interrupt is active if its associated enable and pending bits are set. If GIE = 1 and an interrupt is active, then the processor will be interrupted as soon as it is ready to start executing an instruction except if the above conditions happen during the Software Trap service routine. This exception is described in the Software Trap sub-section.

The interruption process is accomplished with the INTR instruction (opcode 00), which is jammed inside the Instruction Register and replaces the opcode about to be executed. The following steps are performed for every interrupt:

- 1. The GIE (Global Interrupt Enable) bit is reset.
- 2. The address of the instruction about to be executed is pushed into the stack.
- 3. The PC (Program Counter) branches to address 00FF. This procedure takes 7 t_c cycles to execute.

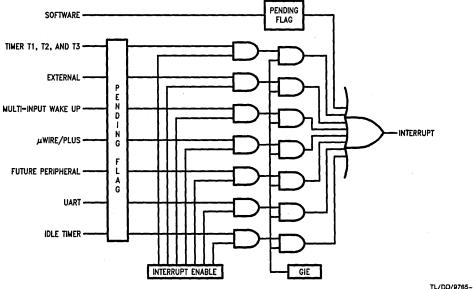


FIGURE 16. COP888CG Interrupt Block Diagram

TL/DD/9765-22

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Interru	pts ((Continued)
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Arbitration Ranking	Source	Description	Vector Address HI-Low Byte
(1) Highest	Software	INTR Instruction	0yFE-0yFF
	Reserved	for Future Use	0yFC-0yFD
(2)	External	Pin G0 Edge	0yFA-0yFB
(3)	Timer T0	Underflow	0yF8-0yF9
(4)	Timer T1	T1A/Underflow	0yF6-0yF7
(5)	Timer T1	T1B	0yF4-0yF5
(6)	MICROWIRE/PLUS	BUSY Goes Low	0yF2-0yF3
	Reserved	for Future Use	0yF0-0yF1
(7)	UART	Receive	0yEE-0yEF
(8)	UART	Transmit	0yEC-0yED
(9)	Timer T2	T2A/Underflow	0yEA-0yEB
(10)	Timer T2	T2B	0yE8-0yE9
(11)	Timer T3	T3A/Underflow	0yE6-0yE7
(12)	Timer T3	ТЗВ	0yE4-0yE5
(13)	Port L/Wakeup	Port L Edge	0yE2-0yE3
(14) Lowest	Default	VIS Instr. Execution without Any Interrupts	0yE0-0yE1

y is VIS page, y ≠ 0.

At this time, since GIE = 0, other maskable interrupts are disabled. The user is now free to do whatever context switching is required by saving the context of the machine in the stack with PUSH instructions. The user would then program a VIS (Vector Interrupt Select) instruction in order to branch to the interrupt service routine of the highest priority interrupt enabled and pending at the time of the VIS. Note that this is not necessarily the interrupt that caused the branch to address location 00FF Hex prior to the context switching.

Thus, if an interrupt with a higher rank than the one which caused the interruption becomes active before the decision of which interrupt to service is made by the VIS, then the interrupt with the higher rank will override any lower ones and will be acknowledged. The lower priority interrupt(s) are still pending, however, and will cause another interrupt immediately following the completion of the interrupt service routine associated with the higher priority interrupt just serviced. This lower priority interrupt will occur immediately following the RETI (Return from Interrupt) instruction at the end of the interrupt service routine just completed.

Inside the interrupt service routine, the associated pending bit has to be cleared by software. The RETI (Return from Interrupt) instruction at the end of the interrupt service routine will set the GIE (Global Interrupt Enable) bit, allowing the processor to be interrupted again if another interrupt is active and pending.

The VIS instruction looks at all the active interrupts at the time it is executed and performs an indirect jump to the beginning of the service routine of the one with the highest rank.

The addresses of the different interrupt service routines, called vectors, are chosen by the user and stored in ROM in a table starting at 01E0 (assuming that VIS is located between 00FF and 01DF). The vectors are 15-bit wide and therefore occupy 2 ROM locations.

VIS and the vector table must be located in the same 256byte block (0y00 to 0yFF) except if VIS is located at the last address of a block. In this case, the table must be in the next block. The vector table cannot be inserted in the first 256-byte block ($y \neq 0$).

The vector of the maskable interrupt with the lowest rank is located at 0yE0 (Hi-Order byte) and 0yE1 (Lo-Order byte) and so forth in increasing rank number. The vector of the maskable interrupt with the highest rank is located at 0yFA (Hi-Order byte) and 0yFB (Lo-Order byte).

The Software Trap has the highest rank and its vector is located at 0yFE and 0yFF.

If, by accident, a VIS gets executed and no interrupt is active, then the PC (Program Counter) will branch to a vector located at 0yE0-0yE1. This vector can point to the Software Trap (ST) interrupt service routine, or to another special service routine as desired.

Figure 16 shows the COP888CG Interrupt block diagram.

SOFTWARE TRAP

The Software Trap (ST) is a special kind of non-maskable interrupt which occurs when the INTR instruction (used to acknowledge interrupts) is fetched from ROM and placed inside the instruction register. This may happen when the PC is pointing beyond the available ROM address space or when the stack is over-popped.

Interrupts (Continued)

When an ST occurs, the user can re-initialize the stack pointer and do a recovery procedure (similar to reset, but not necessarily containing all of the same initialization procedures) before restarting.

The occurrence of an ST is latched into the ST pending bit. The GIE bit is not affected and the ST pending bit (not accessible by the user) is used to inhibit other interrupts and to direct the program to the ST service routine with the VIS Instruction. The RPND instruction is used to clear the software interrupt pending bit. This pending bit is also cleared on reset.

The ST has the highest rank among all interrupts.

Nothing (except another ST) can interrupt an ST being serviced.

WatchDog

The COP888CG contains a WatchDog and clock monitor. The WatchDog is designed to detect the user program getting stuck in infinite loops resulting in loss of program control or "runaway" programs. The Clock Monitor is used to detect the absence of a clock or a very slow clock below a specified rate on the CKI pin.

The WatchDog consists of two independent logic blocks: WD UPPER and WD LOWER. WD UPPER establishes the upper limit on the service window and WD LOWER defines the lower limit of the service window.

Servicing the WatchDog consists of writing a specific value to a WatchDog Service Register named WDSVR which is memory mapped in the RAM. This value is composed of three fields, consisting of a 2-bit Window Select, a 5-bit Key Data field, and the 1-bit Clock Monitor Select field. Table III shows the WDSVR register.

The lower limit of the service window is fixed at 2048 instruction cycles. Bits 7 and 6 of the WDSVR register allow the user to pick an upper limit of the service window.

Table IV shows the four possible combinations of lower and upper limits for the WatchDog service window. This flexibility in choosing the WatchDog service window prevents any undue burden on the user software.

Bits 5, 4, 3, 2 and 1 of the WDSVR register represent the 5bit Key Data field. The key data is fixed at 01100. Bit 0 of the WDSVR Register is the Clock Monitor Select bit.

TABLE III. Wat	chdog Service	e Register (WDSVR)
----------------	---------------	--------------------

	dow ect		Key Data			Clock Monitor	
х	х	0	1	1	0	0	Y
7 ·	6	5	4	3	2	1	0

TABLE IV. Watchdog Service Window Select

WDSVR Bit 7	WDSVR Bit 6	Service Window (Lower-Upper Limits)
0	0	2k-8k t _c Cycles
0	1	2k-16k t _c Cycles
1	0	2k-32k t _c Cycles
1	1	2k-64k t _c Cycles

Clock Monitor

The Clock Monitor aboard the COP888CG can be selected or deselected under program control. The Clock Monitor is guaranteed not to reject the clock if the instruction cycle clock ($1/t_c$) is greater or equal to 10 kHz. This equates to a clock input rate on CKI of greater or equal to 100 kHz.

WatchDog Operation

The WatchDog and Clock Monitor are disabled during reset. The COP888CG comes out of reset with the WatchDog armed, the WatchDog Window Select bits (bits 6, 7 of the WDSVR Register) set, and the Clock Monitor bit (bit 0 of the WDSVR Register) enabled. Thus, a Clock Monitor error will occur after coming out of reset, if the instruction cycle clock frequency has not reached a minimum specified value, including the case where the oscillator fails to start.

The WDSVR register can be written to only once after reset and the key data (bits 5 through 1 of the WDSVR Register) must match to be a valid write. This write to the WDSVR register involves two irrevocable choices: (i) the selection of the WatchDog service window (ii) enabling or disabling of the Clock Monitor. Hence, the first write to WDSVR Register involves selecting or deselecting the Clock Monitor, select the WatchDog service window and match the WatchDog key data. Subsequent writes to the WDSVR register will compare the value being written by the user to the Watch-Dog service window value and the key data (bits 7 through 1) in the WDSVR Register. Table V shows the sequence of events that can occur.

The user must service the WatchDog at least once before the upper limit of the service window expires. The Watch-Dog may not be serviced more than once in every lower limit of the service window. The user may service the WatchDog as many times as wished in the time period between the lower and upper limits of the service window. The first write to the WDSVR Register is also counted as a WatchDog service.

The WatchDog has an output pin associated with it. This is the WDOUT pin, on pin 1 of the port G. WDOUT is active low. The WDOUT pin is in the high impedance state in the inactive state. Upon triggering the WatchDog, the logic will pull the WDOUT (G1) pin low for an additional 16 t_c-32 t_c cycles after the signal level on WDOUT pin goes below the lower Schmitt trigger threshold. After this delay, the COP888CG will stop forcing the WDOUT output low.

The WatchDog service window will restart when the WDOUT pin goes high. It is recommended that the user tie the WDOUT pin back to V_{CC} through a resistor in order to pull WDOUT high.

A WatchDog service while the WDOUT signal is active will be ignored. The state of the WDOUT pin is not guaranteed on reset, but if it powers up low then the WatchDog will time out and WDOUT will enter high impedance state.

The Clock Monitor forces the G1 pin low upon detecting a clock frequency error. The Clock Monitor error will continue until the clock frequency has reached the minimum specified value, after which the G1 output will enter the high impedance TRI-STATE mode following 16 t_c -32 t_c clock cycles. The Clock Monitor generates a continual Clock Monitor error if the oscillator fails to start, or fails to reach the minimum specified frequency. The specification for the Clock Monitor is as follows:

1/t_c > 10 kHz-No clock rejection.

1/t_c < 10 Hz—Guaranteed clock rejection.

Detection of Illegal Conditions

The COP888CG can detect various illegal conditions resulting from coding errors, transient noise, power supply voltage drops, runaway programs, etc.

Reading of undefined ROM gets zeros. The opcode for software interrupt is zero. If the program fetches instructions from undefined ROM, this will force a software interrupt, thus signaling that an illegal condition has occurred.

The subroutine stack grows down for each call (jump to subroutine), interrupt, or PUSH, and grows up for each return or POP. The stack pointer is initialized to RAM location 06F Hex during reset. Consequently, if there are more returns than calls, the stack pointer will point to addresses 070 and 071 Hex (which are undefined RAM). Undefined RAM from addresses 070 to 07F (Segment 0), 140 to 17F (Segment 1), and all other segments (i.e., Segments 3 ... etc.) is read as all 1's, which in turn will cause the program to return to address 7FFF Hex. This is an undefined ROM location and the instruction fetched (all 0's) from this location will generate a software interrupt signaling an illegal condition.

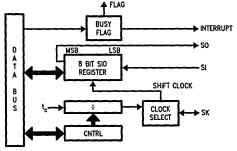
Thus, the chip can detect the following illegal conditions:

- a. Executing from undefined ROM
- b. Over "POP"ing the stack by having more returns than calls.

When the software interrupt occurs, the user can re-initialize the stack pointer and do a recovery procedure before restarting (this recovery program is probably similar to that following reset, but might not contain the same program initialization procedures). The recovery program should reset the software interrupt pending bit using the RPND instruction.

MICROWIRE/PLUS

MICROWIRE/PLUS is a serial synchronous communications interface. The MICROWIRE/PLUS capability enables the COP888CG to interface with any of National Semiconductor's MICROWIRE peripherals (i.e. A/D converters, display drivers, E2PROMs etc.) and with other microcontrollers which support the MICROWIRE interface. It consists of an 8-bit serial shift register (SIO) with serial data input (SI), serial data output (SO) and serial shift clock (SK). Figure 13 shows a block diagram of the MICROWIRE/PLUS logic.



TL/DD/9765-23

FIGURE 17. MICROWIRE/PLUS Block Diagram

The shift clock can be selected from either an internal source or an external source. Operating the MICROWIRE/ PLUS arrangement with the internal clock source is called the Master mode of operation. Similarly, operating the MICROWIRE/PLUS arrangement with an external shift clock is called the Slave mode of operation.

The CNTRL register is used to configure and control the MICROWIRE/PLUS mode. To use the MICROWIRE/PLUS, the MSEL bit in the CNTRL register is set to one. In the master mode, the SK clock rate is selected by the two bits, SL0 and SL1, in the CNTRL register. Table VI details the different clock rates that may be selected.

Key Data	Window Data	Clock Monitor	Action
Match	Match	Match	Valid Service: Restart Service Window
Don't Care	Mismatch	Don't Care	Error: Generate WatchDog Output
Mismatch	Don't Care	Don't Care	Error: Generate WatchDog Output
Don't Care	Don't Care	Mismatch	Error: Generate WatchDog Output

TABLE V. Watchdog Service Actions

TABLE VI. MICROWIRE/PLUS Master Mode Clock Select

SL1	SL0	SK
0	0	$2 \times t_c$
0	1	4×t _c
1	x	$8 \times t_c$

Where t_c is the instruction cycle clock

MICROWIRE/PLUS (Continued)

MICROWIRE/PLUS OPERATION

Setting the BUSY bit in the PSW register causes the MI-CROWIRE/PLUS to start shifting the data. It gets reset when eight data bits have been shifted. The user may reset the BUSY bit by software to allow less than 8 bits to shift. If enabled, an interrupt is generated when eight data bits have been shifted. The COP888CG may enter the MICROWIRE/ PLUS mode either as a Master or as a Slave. *Figure 14* shows how two COP888CG microcontrollers and several peripherals may be interconnected using the MICROWIRE/ PLUS arrangements.

Warning:

The SIO register should only be loaded when the SK clock is low. Loading the SIO register while the SK clock is high will result in undefined data in the SIO register. SK clock is normally low when not shifting.

Setting the BUSY flag when the input SK clock is high in the MICROWIRE/PLUS slave mode may cause the current SK clock forthe SIO shift register to be narrow. For safety, the BUSY flag should only be set when the input SK clock is low.

MICROWIRE/PLUS Master Mode Operation

In the MICROWIRE/PLUS Master mode of operation the shift clock (SK) is generated internally by the COP888CG. The MICROWIRE Master always initiates all data exchanges. The MSEL bit in the CNTRL register must be set to enable the SO and SK functions onto the G Port. The SO and SK pins must also be selected as outputs by setting appropriate bits in the Port G configuration register. Table VII summarizes the bit settings required for Master mode of operation.

MICROWIRE/PLUS Slave Mode Operation

In the MICROWIRE/PLUS Slave mode of operation the SK clock is generated by an external source. Setting the MSEL bit in the CNTRL register enables the SO and SK functions onto the G Port. The SK pin must be selected as an input and the SO pin is selected as an output pin by setting and resetting the appropriate bit in the Port G configuration register. Table VII summarizes the settings required to enter the Slave mode of operation.

The user must set the BUSY flag immediately upon entering the Slave mode. This will ensure that all data bits sent by the Master will be shifted properly. After eight clock pulses the BUSY flag will be cleared and the sequence may be repeated.

Alternate SK Phase Operation

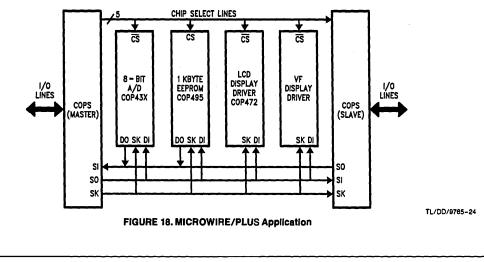
The COP888CG allows either the normal SK clock or an alternate phase SK clock to shift data in and out of the SIO register. In both the modes the SK is normally low. In the normal mode data is shifted in on the rising edge of the SK clock and the data is shifted out on the falling edge of the SK clock. The SIO register is shifted on each falling edge of the SK clock. In the alternate SK phase operation, data is shifted in on the falling edge of the SK clock and shifted out on the rising edge of the SK clock.

A control flag, SKSEL, allows either the normal SK clock or the alternate SK clock to be selected. Resetting SKSEL causes the MICROWIRE/PLUS logic to be clocked from the normal SK signal. Setting the SKSEL flag selects the alternate SK clock. The SKSEL is mapped into the G6 configuration bit. The SKSEL flag will power up in the reset condition, selecting the normal SK signal.

TABLE VII

This table assumes that the control flag MSEL is set.

G4 (SO) Config. Bit	G5 (SK) Config. Bit	G4 Fun.	G5 Fun.	Operation
1	1	SO	int. SK	MICROWIRE/PLUS Master
0	1	TRI- STATE	lnt. SK	MICROWIRE/PLUS Master
1	0	SO	Ext. SK	MICROWIRE/PLUS Slave
0	0	TRI- STATE	Ext. SK	MICROWIRE/PLUS Slave



Memory Map

All RAM, ports and registers (except A and PC) are mapped into data memory address space.

Address S/ADD REG	Contents
0000 to 006F	On-Chip RAM bytes (112 bytes)
0070 to 007F xx80 to xxAF	Unused RAM Address Space (Reads As All Ones) Unused RAM Address Space (Reads
	Undefined Data)
xxB0	Timer T3 Lower Byte
XXB1	Timer T3 Upper Byte
xxB2	Timer T3 Autoload Register T3RA Lower Byte
xxB3	Timer T3 Autoload Register T3RA Upper Byte
xxB4	Timer T3 Autoload Register T3RB Lower Byte
xxB5	Timer T3 Autoload Register T3RB Upper Byte
xxB6	Timer T3 Control Register
xxB7	Comparator Select Register (CMPSL)
xxB8	UART Transmit Buffer (TBUF)
xxB9	UART Receive Buffer (RBUF)
ххВА	UART Control and Status Register (ENU)
XXBB	UART Receive Control and Status Register (ENUR)
XXBC	UART Interrupt and Clock Source Register (ENUI)
ххBD	UART Baud Register (BAUD)
XXBE	UART Prescale Select Register (PSR)
xxBF	Reserved for UART
xxC0	Timer T2 Lower Byte
xxC1	Timer T2 Upper Byte
xxC2	Timer T2 Autoload Register T2RA Lower Byte
xxC3	Timer T2 Autoload Register T2RA Upper Byte
xxC4	Timer T2 Autoload Register T2RB Lower Byte
xxC5	Timer T2 Autoload Register T2RB Upper Byte
xxC6	Timer T2 Control Register
xxC7	WatchDog Service Register (Reg:WDSVR)
xxC8	MIWU Edge Select Register (Reg:WKEDG)
xxC9	MIWU Enable Register (Reg:WKEN)
XXCA	MIWU Pending Register (Reg:WKPND)
xxCB	Reserved
XXCC	Reserved
xxCD to xxCF	Reserved

ta memory address space.					
Address S/ADD REG	Contents				
xxD0	Port L Data Register				
xxD1	Port L Configuration Register				
xxD2	Port L Input Pins (Read Only)				
xxD3	Reserved for Port L				
xxD4	Port G Data Register				
xxD5	Port G Configuration Register				
xxD6	Port G Input Pins (Read Only)				
xxD7	Port I Input Pins (Read Only)				
xxD8	Port C Data Register				
xxD9	Port C Configuration Register				
xxDA	Port C Input Pins (Read Only)				
xxDB	Reserved for Port C				
xxDC	Port D				
xxDD to DF	Reserved for Port D				
xxE0 to xxE5	Reserved for EE Control Registers				
xxE6	Timer T1 Autoload Register T1RB				
	Lower Byte				
xxE7	Timer T1 Autoload Register T1RB				
	Upper Byte				
XXE8	ICNTRL Register				
xxE9	MICROWIRE/PLUS Shift Register				
ххEA	Timer T1 Lower Byte				
XXEB	Timer T1 Upper Byte				
xxEC	Timer T1 Autoload Register T1RA				
	Lower Byte				
xxED	Timer T1 Autoload Register T1RA				
	Upper Byte				
XXEE	CNTRL Control Register				
xxEF	PSW Register				
xxF0 to FB	On-Chip RAM Mapped as Registers				
xxFC	X Register				
xxFD	SP Register				
xxFE	B Register				
xxFF	S Register				
0100-013F	On-Chip RAM Bytes (64 bytes)				

Reading memory locations 0070H-007FH (Segment 0) will return all ones. Reading unused memory locations 0080H-00AFH (Segment 0) will return undefined data. Reading unused memory locations 0140-017F (Segment 1) will return all ones. Reading memory locations from other Segments (i.e., Segment 2, Segment 3, ... etc.) will return all ones.

Addressing Modes

The COP888CG has ten addressing modes, six for operand addressing and four for transfer of control.

OPERAND ADDRESSING MODES

Register Indirect

This is the "normal" addressing mode for the COP888CG. The operand is the data memory addressed by the B pointer or X pointer.

Register Indirect (with auto post increment or decrement of pointer)

This addressing mode is used with the LD and X instructions. The operand is the data memory addressed by the B pointer or X pointer. This is a register indirect mode that automatically post increments or decrements the B or X register after executing the instruction.

Direct

The instruction contains an 8-bit address field that directly points to the data memory for the operand.

Immediate

The instruction contains an 8-bit immediate field as the operand.

Short Immediate

This addressing mode is used with the LBI instruction. The instruction contains a 4-bit immediate field as the operand.

Indirect

This addressing mode is used with the LAID instruction. The contents of the accumulator are used as a partial address (lower 8 bits of PC) for accessing a data operand from the program memory.

TRANSFER OF CONTROL ADDRESSING MODES

Relative

This mode is used for the JP instruction, with the instruction field being added to the program counter to get the new program location. JP has a range from -31 to +32 to allow a 1-byte relative jump (JP + 1 is implemented by a NOP instruction). There are no "pages" when using JP, since all 15 bits of PC are used.

Absolute

This mode is used with the JMP and JSR instructions, with the instruction field of 12 bits replacing the lower 12 bits of the program counter (PC). This allows jumping to any location in the current 4k program memory segment.

Absolute Long

This mode is used with the JMPL and JSRL instructions, with the instruction field of 15 bits replacing the entire 15 bits of the program counter (PC). This allows jumping to any location in the current 4k program memory space.

Indirect

This mode is used with the JID instruction. The contents of the accumulator are used as a partial address (lower 8 bits of PC) for accessing a location in the program memory. The contents of this program memory location serve as a partial address (lower 8 bits of PC) for the jump to the next instruction.

Note: The VIS is a special case of the Indirect Transfer of Control addresslng mode, where the double byte vector associated with the interrupt is transferred from adjacent addresses in the program memory into the program counter (PC) in order to jump to the associated interrupt service routine.

Instruction Set

Register and Symbol Definition

Registers					
Α	8-Bit Accumulator Register				
в	8-Bit Address Register				
х	8-Bit Address Register				
SP	8-Bit Stack Pointer Register				
PC .	15-Bit Program Counter Register				
PU	Upper 7 Bits of PC				
PL	Lower 8 Bits of PC				
C	1 Bit of PSW Register for Carry				
HC	1 Bit of PSW Register for Half Carry				
GIE	1 Bit of PSW Register for Global				
	Interrupt Enable				
VU	Interrupt Vector Upper Byte				
VL	Interrupt Vector Lower Byte				

Symbols				
(B)	Memory Indirectly Addressed by B Register			
[X]	Memory Indirectly Addressed by X Register			
MD	Direct Addressed Memory			
Mem	Direct Addressed Memory or [B]			
Memi	Direct Addressed Memory or [B] or Immediate Data			
lmm	8-Bit Immediate Data			
Reg	Register Mernory: Addresses F0 to FF (Includes B, X and SP)			
Bit	Bit Number (0 to 7)			
←	Loaded with			
\longleftrightarrow	Exchanged with			

Instruction Set (Continued)

INSTRUCTION SET

$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$
SUBCA,MemlSubtract with Carry $HC \leftarrow Half Carry$ ANDA,MemlLogical AND $A \leftarrow A$ and Meml $C \leftarrow Carry$ ANDSZA,ImmLogical AND Immed., Skip if ZeroSkip next if (A and Imm) = 0ORA,MemlLogical EXclusive OR $A \leftarrow A$ and MemlXORA,MemlLogical EXclusive OR $A \leftarrow A$ or MemlIFEQMD,ImmIF EQualCompare A and Meml, Do next if MD = ImmIFEQA,MemlIF EQualCompare A and Meml, Do next if A = MemlIFEQA,MemlIF Rot EqualCompare A and Meml, Do next if A = MemlIFSTA,MemlIF StrathCompare A and Meml, Do next if A > MemlIFSNE#If B Not EqualDo next if lower 4 bits of B \neq ImmDRSZRegDecrement Reg., Skip if ZeroReg \leftarrow Reg -1 , Skip if Reg $= 0$ SBIT#,MemSet BITI to bit, Mem (bit $= 0$ to 7 immediate)RBIT#,MemReset BITIf bit in A or Mem is true do next instructionRPNDReset PeNDing FlagReset Software Interrupt Pending FlagXA,[X]LoaD A with Memory $A \leftarrow Meml$ LDA,[X]LoaD A with Memory [X] $A \leftarrow IX]$ LDB,ImmLoaD A with Memory Immed $A \leftarrow Imm$ LDReg,ImmLoaD Register Memory ImmedReg \leftarrow ImmLDReg,ImmLoaD Register Memory ImmedReg \leftarrow ImmLDReg,ImmLoaD Register Memory ImmedReg \leftarrow Imm
SUBCA,MemlSubtract with Carry $HC \leftarrow Half Carry$ ANDA,MemlLogical AND $A \leftarrow A$ and Meml $C \leftarrow Carry$ ANDSZA,ImmLogical AND Immed., Skip if ZeroSkip next if (A and Imm) = 0ORA,MemlLogical CR $A \leftarrow A$ and MemlXORA,MemlLogical EXclusive OR $A \leftarrow A$ or MemlIFEQMD,ImmIF EQualCompare A ond Meml, Do next if MD = ImmIFEQA,MemlIF EQualCompare A and Meml, Do next if A = MemlIFERA,MemlIF Rot EqualCompare A and Meml, Do next if A = MemlIFRA,MemlIF Rotaer ThanCompare A and Meml, Do next if A > MemlIFBNE#If B Not EqualDo next if lower 4 bits of B \neq ImmDRSZRegDecrement Reg., Skip if ZeroReg \leftarrow Reg -1 , Skip if Reg $= 0$ SBIT#,MemSet BITI to bit, Mem (bit = 0 to 7 immediate)RBIT#,MemIF BITIf BITRPNDReset PeNDing FlagReset Software Interrupt Pending FlagXA,[X]LoaD A with Memory $A \leftarrow Meml$ LDA,[X]LoaD A with Memory [X] $A \leftarrow I[X]$ LDB,ImmLoaD A with Memory Immed $A \leftarrow Imm$ LDReg,ImmLoaD Register Memory Immed. $Reg \leftarrow Imm$ XA, [B ± 1]EXchange A with Memory [B] $A \leftarrow Meml$
SUBCA,MemlSubtract with Carry $A \leftarrow A - Meml + C, C \leftarrow Carry + HC \leftarrow Half Hard + Hard$
ANDA,MemlLogical AND $HC \leftarrow Half Carry$ ANDSZA,ImmLogical AND Immed., Skip if ZeroSkip next if (A and Imm) = 0ORA,MemlLogical EXclusive ORSkip next if (A and Imm) = 0IFEQMD,ImmIF EQualCompare A or MemlIFEQMD,ImmIF EQualCompare A and Meml, Do next if MD = ImmIFEQA,MemlIF EQualCompare A and Meml, Do next if A = MemlIFEQA,MemlIF Reater ThanCompare A and Meml, Do next if A = MemlIFBNE#If B Not EqualDo next if Iower 4 bits of B ≠ ImmDRSZRegDecrement Reg., Skip if ZeroReg ← Reg - 1, Skip if Reg = 0SBIT#,MemSet BITIt bit in A or Mem is true do next instructionRBIT#,MemReset BITIf bit in A or Mem is true do next instructionRPNDReset PeNDing FlagA ← MemiXA,[X]LoaD A with MemoryA ← MemiLDA,[X]LoaD A with Memory [X]A ← [X]LDB,ImmLoaD A with Memory [X]A ← [X]LDReg,ImmLoaD Register Memory Immed.B ← ImmLDReg,ImmLoaD Register Memory Immed.B ← ImmLDReg,ImmLoaD Register Memory Immed.A ← EXA, [B ±]EXchange A with Memory [B]A ← [B], (B ← B ± 1)
ANDA,MemlLogical ANDA \leftarrow A and MemlANDSZA,ImmLogical AND Immed., Skip if ZeroSkip next if (A and Imm) = 0ORA,MemlLogical ORA \leftarrow A or MemlXORA,MemlLogical EXclusive ORA \leftarrow A or MemlIFEQMD,ImmIF EQualCompare M and Meml, Do next if MD = ImmIFEQA,MemlIF EQualCompare A and Meml, Do next if A = MemlIFEQA,MemlIF Rot EqualCompare A and Meml, Do next if A = MemlIFBNE#If B Not EqualCompare A and Meml, Do next if A > MemlIFBNE#If B Not EqualDo next if Iower 4 bits of B ≠ ImmDRSZRegDecrement Reg., Skip if ZeroReg \leftarrow Reg - 1, Skip if Reg = 0SBIT#,MemReset BITI to bit, Mem (bit = 0 to 7 immediate)IFBIT#,MemIF BITIf BITIFBIT#,MemIF BITIFBIT#,MemEXchange A with MemoryXA,[X]EXchange A with Memory [X]AAAXA,[X]LoaD A with Memory [X]LDA,ImmLoaD A with Memory [X]LDB,ImmLoaD A with Memory [X]LDB,ImmLoaD A with Memory Immed.LDReg,ImmLoaD Register Memory Immed.LDReg,Imm<
ANDSZA,ImmLogical AND Immed., Skip if ZeroSkip next if (A and Imm) = 0ORA,MemlLogical OR $A \leftarrow A$ or MemlXORA,MemlLogical EXclusive OR $A \leftarrow A$ or MemlIFEQMD,ImmIF EQualCompare MD and Imm, Do next if MD = ImmIFEQA,MemlIF EQualCompare A and Meml, Do next if A = MemlIFNEA,MemlIF Greater ThanCompare A and Meml, Do next if A > MemlIFBNE#If B Not EqualDo next if Iower 4 bits of B ≠ ImmDRSZRegDecrement Reg., Skip if ZeroReg \leftarrow Reg - 1, Skip if Reg = 0SBIT#,MemSet BIT1 to bit, Mem (bit = 0 to 7 immediate)IFBIT#,MemIF BITIf B ITIFBIT#,MemEXchange A with MemoryA ← → MemXA,[X]EXchange A with Memory [X]A ←→ [X]LDA,MemlLoaD A with Memory [X]A ←→ [X]LDB,ImmLoaD A with Memory [X]A ←→ [X]LDB,ImmLoaD A with Memory Immed.B ← ImmLDReg,ImmLoaD Register Memory Immed.Reg ← ImmXA, [B ±]EXchange A with Memory Immed.Reg ← Imm
ORA,MemlLogical OR $A \leftarrow A \text{ or Meml}$ XORA,MemlLogical EXclusive OR $A \leftarrow A \text{ or Meml}$ IFEQMD,ImmIF EQualCompare MD and Imm, Do next if MD = ImmIFEQA,MemlIF EQualCompare A and Meml, Do next if A = MemlIFNEA,MemlIF Not EqualCompare A and Meml, Do next if A > MemlIFBNE#If B Not EqualDo next if Iower 4 bits of B ≠ ImmDRSZRegDecrement Reg., Skip if ZeroReg ← Reg - 1, Skip if Reg = 0SBIT#,MemSet BIT1 to bit, Mem (bit = 0 to 7 immediate)IFBIT#,MemReset BITIf b bit in A or Mem is true do next instructionRPNDReset PeNDing FlagReset Software Interrupt Pending FlagXA,[X]LoaD A with Memory (X] $A \leftarrow Meml$ LDB,ImmLoaD A with Memory (X] $A \leftarrow Meml$ LDB,ImmLoaD A with Memory Immed. $B \leftarrow Imm$ LDReg,ImmLoaD Register Memory Immed. $A \leftarrow F[R]$, $(B \leftarrow B \pm 1)$
XORA,MemlLogical EXclusive OR $A \leftarrow A \text{ xor Meml}$ IFEQMD,ImmIF EQualCompare A and Meml, Do next if MD = ImmIFEQA,MemlIF EQualCompare A and Meml, Do next if A = MemlIFNEA,MemlIF Not EqualCompare A and Meml, Do next if A = MemlIFBNE#If B Not EqualCompare A and Meml, Do next if A > MemlDRSZRegDecrement Reg., Skip if ZeroReg \leftarrow Reg - 1, Skip if Reg = 0SBIT#,MemSet BIT1 to bit, Mem (bit = 0 to 7 immediate)RBIT#,MemIF BITIf b IT in A or Mem is true do next instructionRPNDReset PeNDing FlagReset Software Interrupt Pending FlagXA,[X]EXchange A with MemoryA ← MemlLDA,IXILoaD A with Memory [X]A ← [X]LDB,ImmLoaD A with Memory Immed.B ← ImmLDReg,ImmLoaD Register Memory Immed.Reg ← ImmXA, [B ±]EXchange A with Memory [B]A ←→ [B], (B ← B ± 1)
IFEQMD,ImmIF EQualCompare MD and Imm, Do next if MD = ImmIFEQA,MemIIF EQualCompare A and MemI, Do next if A = MemIIFNEA,MemIIF Not EqualCompare A and MemI, Do next if A = MemIIFNEA,MemIIF Greater ThanCompare A and MemI, Do next if A > MemIIFBNE#If B Not EqualDo next if If A > MemIDRSZRegDecrement Reg., Skip if ZeroReg ← Reg - 1, Skip if Reg = 0SBIT#,MemSet BIT1 to bit, Mem (bit = 0 to 7 immediate)RBIT#,MemIF BITIf B ITIFBIT#,MemIF BITRPNDReset PeNDing FlagReset Software Interrupt Pending FlagXA,IXIEXchange A with Memory [X]A ←→ [X]LDA,MemILoaD A with Memory [X]A ←→ [X]LDB,ImmLoaD A with Memory [X]A ←→ [X]LDB,ImmLoaD B with Immed.B ← ImmLDReg,ImmLoaD Register Memory Immed.Reg ← ImmXA, [B ±]EXchange A with Memory [B]A ←→ [B], (B ← B ± 1)
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SBIT#,MemSet BIT1 to bit, Mem (bit = 0 to 7 immediate)RBIT#,MemReset BIT0 to bit, Mem (bit = 0 to 7 immediate)IFBIT#,MemIF BIT0 to bit, MemRPNDReset PeNDing FlagIf bit in A or Mem is true do next instruction Reset Software Interrupt Pending FlagXA,MemEXchange A with Memory LDA $\leftrightarrow Mem$ LDA,MemILoaD A with Memory [X] LoaD A with Memory [X]A $\leftarrow MemI$ LDB,Imm LoaD B with Immed.B \leftarrow Imm Mem \leftarrow Imm LDLDMem,Imm LoaD B with Immed.B \leftarrow Imm Mem \leftarrow ImmLDReg,ImmLoaD Register Memory Immed.Reg \leftarrow ImmXA, [B ±]Exchange A with Memory [B]A $\leftarrow \Rightarrow$ [B], (B \leftarrow B ± 1)
RBIT# MemReset BIT0 to bit, MemIFBIT# MemIF BITIf bit in A or Mem is true do next instructionRPNDReset PeNDing FlagIf bit in A or Mem is true do next instructionXA,MemEXchange A with MemoryA $\leftrightarrow Mem$ XA,[X]EXchange A with Memory [X]A $\leftrightarrow Mem$ LDA,MemlLoaD A with Memory [X]A $\leftarrow Meml$ LDB,ImmLoaD A with Memory [X]A $\leftarrow IX]$ LDB,ImmLoaD A with Memory [X]A $\leftarrow IX]$ LDB,ImmLoaD A with Memory [X]A $\leftarrow IX]$ LDB,ImmLoaD A with Memory [Memed.B \leftarrow ImmLDReg,ImmLoaD Register Memory Immed.Reg \leftarrow ImmXA, [B ±]EXchange A with Memory [B]A $\leftarrow > [B], (B \leftarrow B \pm 1)$
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$\begin{tabular}{ c c c c c c c c c c c c c c c c c c c$
XA,MemEXchange A with MemoryA \leftrightarrow MemXA,[X]EXchange A with Memory [X]A \leftrightarrow [X]LDA,MemlLoaD A with Memory [X]A \leftarrow [X]LDA,[X]LoaD A with Memory [X]A \leftarrow [X]LDA,[X]LoaD A with Memory [X]A \leftarrow [X]LDB,ImmLoaD B with Immed.B \leftarrow ImmLDMem,ImmLoaD Memory ImmedMem \leftarrow ImmLDReg,ImmLoaD Register Memory Immed.Reg \leftarrow ImmXA, [B \pm]EXchange A with Memory [B]A \leftrightarrow [B], (B \leftarrow B \pm 1)
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$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$
LD Mem,Imm LoaD Memory Immed Mem ← Imm LD Reg,Imm LoaD Register Memory Immed. Reg ← Imm X A, [B ±] EXchange A with Memory [B] A ↔ [B], (B ← B ± 1)
LD Reg,Imm LoaD Register Memory Immed. Reg ← Imm X A, [B ±] EXchange A with Memory [B] A ↔ [B], (B ↔ B ± 1)
XA, $[B \pm]$ EXchange A with Memory $[B]$ A $\leftrightarrow = [B], (B \leftarrow B \pm 1)$
X A, [B \pm] EXchange A with Memory [B] A \leftrightarrow [B], (B \leftarrow B \pm 1)
$[X A, [X \pm]] = Exchange A with Memory [X] A \leftrightarrow [X], (X \leftarrow \pm 1)$
LD A, $(B\pm)$ LoaD A with Memory (B) A \leftarrow (B) , $(B \leftarrow B \pm 1)$
LD A, $[X\pm]$ LoaD A with Memory $[X]$ A \leftarrow $[X], (X \leftarrow X\pm 1)$
LD [B±],Imm LoaD Memory [B] Immed. [B] ← Imm, (B ← B±1)
CLR A CLeaRA A ← 0
INC A INCrement A A + 1
DEC A DECrementA A - 1
LAID Load A InDirect from ROM A ← ROM (PU,A)
DCOR A Decimal CORrect A A ← BCD correction of A (follows ADC, SUBC)
RRC A Rotate A Right thru C $C \rightarrow A7 \rightarrow \rightarrow A0 \rightarrow C$
RLC A Rotate A Left thru C $C \leftarrow A7 \leftarrow \leftarrow A0 \leftarrow C$
SWAP A SWAP nibbles of A $A7A4 \leftrightarrow A3A0$
SC Set C $C \leftarrow 1, HC \leftarrow 1$
RC Reset C $C \leftarrow 0, HC \leftarrow 0$
IFC IF C
IFNC IF Not C If C is not true, do next instruction
POP A POP the stack into A $SP \leftarrow SP + 1, A \leftarrow [SP]$
PUSH A PUSH A onto the stack $[SP] \leftarrow A, SP \leftarrow SP - 1$
VIS Vector to Interrupt Service Routine $PU \leftarrow [VU], PL \leftarrow [VL]$
JMPL Addr. Jump absolute Long PC ← ii (ii = 15 bits, 0 to 32k)
JMP Addr. Jump absolute $PC9 \dots 0 \leftarrow i (i = 12 \text{ bits})$
JP Disp. Jump relative short $PC \leftarrow PC + r (r \text{ is } -31 \text{ to } +32, \text{ except } 1)$
JSRL Addr. Jump SubRoutine Long [SP] ← PL, [SP-1] ← PU,SP-2, PC ← ii
JSR Addr Jump SubRoutine [SP] ← PL, [SP-1] ← PU, SP-2, PC90 <
JID Jump InDirect PL ← ROM (PU,A)
RET RETurn from subroutine SP + 2, PL ← [SP], PU ← [SP-1]
RETSK RETurn and SKip SP + 2, PL ← [SP], PU ← [SP - 1]
RETSKRETurn and Skip $SP + 2, PL \leftarrow [SP], PU \leftarrow [SP-1]$ RETIRETurn from Interrupt $SP + 2, PL \leftarrow [SP], PU \leftarrow [SP-1], GIE \leftarrow 1$ INTRGenerate an Interrupt $[SP] \leftarrow PL, [SP-1] \leftarrow PU, SP-2, PC \leftarrow OFF$
RETSK RETurn and SKip SP + 2, PL ← [SP],PU ← [SP-1] RETI RETurn from Interrupt SP + 2, PL ← [SP],PU ← [SP-1],GIE ← 1

COP888CG

Instruction Execution Time

Most instructions are single byte (with immediate addressing mode instructions taking two bytes).

Most single byte instructions take one cycle time (1 μ s at 10 MHz) to execute.

See the BYTES and CYCLES per INSTRUCTION table for details.

Bytes and Cycles per Instruction

The following table shows the number of bytes and cycles for each instruction in the format of byte/cycle (a cycle is 1 μ s at 10 MHz).

	(B)	Direct	Immed.
ADD	1/1	3/4	2/2
ADC	1/1	3/4	2/2
SUBC	1/1	3/4	2/2
AND	1/1	3/4	2/2
OR	1/1	3/4	2/2
XOR	1/1	3/4	2/2
IFEQ	1/1	3/4	2/2
IFNE	1/1	3/4	2/2
IFGT	1/1	3/4	2/2
IFBNE	1/1		
DRSZ		1/3	
SBIT	1/1	3/4	
RBIT	1/1	3/4	
IFBIT	1/1	3/4	

Instructions Using A & C		Transfer of	••••••
CLRA	1/1	Instruct	lions
INCA	1/1	JMPL	3/4
DECA	1/1	JMP	2/3
LAID ·	1/3	JP	1/3
DCOR	1/1	JSRL	3/5
RRCA	1/1	JSR	2/5
RLCA	1/1	JID	1/3
SWAPA	1/1	VIS	1/5
SC	1/1	RET	1/5
RC	1/1	RETSK	1/5
IFC	1/1	RETI	1/5
IFNC	1/1	INTR	1/7
PUSHA	1/3	NOP	1/1
POPA	1/3	L	4
ANDSZ	2/2		

RPND 1/1

Memory Transfer Instructions

	Register Indirect				Immed.	-	Indirect r. & Decr.	
	(B)	[X]			[B+,B-]	[X+,X-]		
X A,*	1/1	1/3	2/3		1/2	1/3		
LD A,*	1/1	1/3	2/3	2/2	1/2	1/3		
LD B, Imm			-	1/1		}	(IF B < 1	
LD B, Imm				2/2	1		(IF B > 1	
LD Mem, Imm	2/2		3/3		2/2			
LD Reg, Imm			2/3	ļ	ł			
IFEQ MD, Imm			3/3	ļ				

• = > Memory location addressed by B or X or directly.

COP8888CG Opcode Table Upper Nibble Along X-Axis

Lower Nibble Along Y-Axis

F	E	D	C	В	A	9	8	
JP – 15	JP -31	LD 0F0, # i	DRSZ 0F0	RRCA	RC	ADC A,#i	ADC A,[B]	0
JP - 14	JP -30	LD 0F1, # i	DRSZ 0F1	•	SC	SUBC A, #i	SUB A,[B]	1
JP – 13	JP - 29	LD 0F2, # i	DRSZ 0F2	X A, [X+]	X A,[B+]	IFEQ A, #i	IFEQ A,[B]	2
JP - 12	JP - 28	LD 0F3, # i	DRSZ 0F3	X A, [X –]	X A,[B]	IFGT A,#i	IFGT A,[B]	3
JP - 11	JP - 27	LD 0F4, # i	DRSZ 0F4	VIS	LAID	ADD A, #i	ADD A,[B]	4
JP - 10	JP - 26	LD 0F5, # i	DRSZ 0F5	RPND	JID	AND A, #i	AND A,[B]	5
JP -9	JP -25	LD 0F6, # i	DRSZ 0F6	X A,[X]	X A,[B]	XOR A, #i	XOR A,[B]	6
JP -8	JP -24	LD 0F7, # i	DRSZ 0F7	•	•	OR A, #i	OR A,[B]	7
JP -7	JP -23	LD 0F8, # i	DRSZ 0F8	NOP	RLCA	LD A,#i	IFC	8
JP -6	JP - 22	LD 0F9, # i	DRSZ 0F9	IFNE A,[B]	IFEQ Md,#i	IFNE A,#i	IFNC	9
JP - 5	JP -21	LD 0FA, # i	DRSZ OFA	LD A,[X+]	LD A, [B+]	LD [B+],#i	INCA	A
JP -4	JP - 20	LD 0FB, # i	DRSZ 0FB	LD A,[X-]	LD A,[B-)	LD [B-],#i	DECA	В
JP -3	JP - 19	LD 0FC, # i	DRSZ 0FC	LD Md,#i	JMPL	X A,Md	POPA	С
JP -2	JP - 18	LD 0FD, # i	DRSZ 0FD	DIR	JSRL	LD A,Md	RETSK	D
JP - 1	JP - 17	LD OFE, # i	DRSZ OFE	LD A,[X]	LD A,[B]	LD [B],#i	RET	E
JP -0	JP - 16	LD 0FF, # i	DRSZ OFF	•	•	LD B,#i	RETI	F

COP888CG Opcode Table (Continued)

Upper Nibble Along X-Axis

Lower Nibble Along Y-Axis

7	6	5	4	3	2	1	0	
IFBIT 0,[B]	ANDSZ A, #i	LD B,#0F	IFBNE 0	JSR x000-x0FF	JMP x000-x0FF	JP +17	INTR	0
IFBIT 1,[B]	•	LD B,#0E	IFBNE 1	JSR x100-x1FF	JMP x100-x1FF	JP + 18	JP + 2	1
IFBIT 2,[B]	•	LD B,#0D	IFBNE 2	JSR x200-x2FF	JMP x200-x2FF	JP + 19	JP + 3	2
IFBIT 3,[B]	•	LD B,#0C	IFBNE 3	JSR x300-x3FF	JMP x300-x3FF	JP +20	JP + 4	3
IFBIT 4,[B]	CLRA	LD B, #0B	IFBNE 4	JSR x400-x4FF	JMP x400-x4FF	JP +21	JP + 5	4
IFBIT 5,[B]	SWAPA	LD B, #0A	IFBNE 5	JSR x500-x5FF	JMP x500-x5FF	JP + 22	JP + 6	5
IFBIT 6,[B]	DCORA	LD B,#09	IFBNE 6	JSR x600-x6FF	JMP x600-x6FF	JP +23	JP + 7	6
IFBIT 7,[B]	PUSHA	LD B, #08	IFBNE 7	JSR x700-x7FF	JMP x700-x7FF	JP +24	JP + 8	7
SBIT 0,[B]	RBIT 0,[B]	LD B, #07	IFBNE 8	JSR x800-x8FF	JMP x800-x8FF	JP +25	JP + 9	8
SBIT 1,(B]	RBIT 1,[B]	LD B,#06	IFBNE 9	JSR x900-x9FF	JMP x900-x9FF	JP + 26	JP + 10	9
SBIT 2,[B]	RBIT 2,[B]	LD B, #05	IFBNE 0A	JSR xA00-xAFF	JMP xA00-xAFF	JP +27	JP + 11	A
SBIT 3,[B]	RBIT 3,[B]	LD B, #04	IFBNE 0B	JSR xB00-xBFF	JMP xB00-xBFF	JP +28	JP + 12	В
SBIT 4,[B]	RBIT 4,[B]	LD B,#03	IFBNE OC	JSR xC00-xCFF	JMP xC00-xCFF	JP + 29	JP + 13	С
SBIT 5,[B]	RBIT 5,[B]	LD B, #02	IFBNE OD	JSR xD00-xDFF	JMP xD00-xDFF	JP +30	JP + 14	D
SBIT 6,[B]	RBIT 6,[B]	LD B,#01	IFBNE OE	JSR xE00-xEFF	JMP xE00-xEFF	JP +31	JP + 15	E
SBIT 7,[B]	RBIT 7,[B]	LD B,#00	IFBNE OF	JSR xF00-xFFF	JMP xF00-xFFF	JP +32	JP + 16	F

Where,

I is the immediate data

Md is a directly addressed memory location

* is an unused opcode

Note: The opcode 60 Hex is also the opcode for IFBIT #I,A

COP888CG

Mask Options

The COP888CG mask programmable options are shown below. The options are programmed at the same time as the ROM pattern submission.

OPTION 1: CLOCK CONFIGURATION

- = 1 Crystal Oscillator (CKI/10)
 G7 (CKO) is clock generator
 output to crystal/resonator
 CKI is the clock input
 - Single-pin RC controlled oscillator (CKI/10) G7 is available as a HALT restart and/or general purpose input

OPTION 2: HALT

= 2

- = 1 Enable HALT mode
- = 2 Disable HALT mode

OPTION 3: COP888CG BONDING

- = 1 44-Pin PCC
- = 2 40-Pin DIP
- = 3 28-Pin PCC
- = 4 28-Pin DIP

The chip can be driven by a clock input on the CKI input pin which can be between DC and 10 MHz. The CKO output clock is on pin G7 (if clock option-1 has been selected). The CKI input frequency is divided down by 10 to produce the instruction cycle clock $(1/t_c)$.

Development Support

MOLE DEVELOPMENT SYSTEM

The MOLE (Microcomputer On Line Emulator) is a low cost development system and emulator for all microcontroller products. These include COPs™ microcontrollers and the HPC family of products. The MOLE consists of a BRAIN Board, Personality Board and optional host software.

The purpose of the MOLE is to provide the user with a tool to write and assemble code, emulate code for the target microcontroller and assist in both software and hardware debugging of the system.

It is a self contained computer with its own firmware which provides for all system operation, emulation control, communication, PROM programming and diagnostic operations.

It contains three serial ports to optionally connect to a terminal, a host system, a printer or a modem, or to connect to other MOLEs in a multi-MOLE environment.

MOLE can be used in either a stand alone mode or in conjunction with a selected host system using PC-DOS communicating via a RS-232 port.

How to Order

To order a complete development package, select the section for the microcontroller to be developed and order the parts listed.

Development Tools Selection Table

Microcontroller	Order Part Number	Description	Includes	Manual Number
	MOLE-BRAIN	Brain Board	Brain Board Users Manual	420408188-001
	MOLE-COP8-PB2	Personality Board	COP888 Personality Board Users Manual	420420084-001
COP888	MOLE-COP8-IBM	Assembler Software for IBM	COP800 Software Users Manual and Software Disk PC-DOS Communications Software Users Manual	424410527-001 420040416-001
	420411060-001	Programmer's Manual		420411060-01

Development Support (Continued) DIAL-A-HELPER

Dial-A-Helper is a service provided by the Microcontroller Applications group. The Dial-A-Helper is an Electronic Bulletin Board Information system and additionally, provides the capability of remotely accessing the MOLE development system at a customer site.

INFORMATION SYSTEM

The Dial-A-Helper system provides access to an automated information storage and retrieval system that may be accessed over standard dial-up telephone lines 24 hours a day. The system capabilities include a MESSAGE SECTION (electronic mail) for communications to and from the Microcontroller Applications Group and a FILE SECTION which consists of several file areas where valuable application software and utilities could be found. The minimum requirement for accessing the Dial-A-Helper is a Hayes compatible modem. If the user has a PC with a communications package then files from the FILE SECTION can be down loaded to disk for later use.

ORDER P/N: MOLE-DIAL-A-HLP

Information System Package contains: Dial-A-Helper Users Manual Public Domain Communications Software

FACTORY APPLICATIONS SUPPORT

Dial-A-Helper also provides immediate factor applications support. If a user is having difficulty in operating a MOLE, he can leave messages on our electronic bulletin board, which we will respond to, or under extraordinary circumstances he can arrange for us to actually take control of his system via modem for debugging purposes.

